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FY17 ASC P&EM L2 Milestone 6016: Tri-Lab Co-Design Milestone: In-Depth Performance Portability Analysis of Improved Integrated Codes on Advanced Architecture

Executive Summary
Author: Robert Hoekstra, 09/22/2017

Introduction
This milestone is a tri-lab deliverable supporting ongoing Co-Design efforts impacting applications in the Integrated Codes (IC) program element Advanced Technology Development and Mitigation (ATDM) program element. In FY14, the trilabs looked at porting proxy application to technologies of interest for ATS procurements. In FY15, a milestone was completed evaluating proxy applications in multiple programming models and in FY16, a milestone was completed focusing on the migration of lessons learned back into production code development. This year, the co-design milestone focuses on extracting the knowledge gained and/or code revisions back into production applications.

Milestone Description
This milestone will exercise the IC or ATDM codes using new and emerging programming models, to demonstrate portability and analyze performance characteristics on the latest advanced architectures expected to become widely available in FY17. Performance analysis tools will be used to do an in-depth analysis of those same or similar production codes on available ASC AT platforms and test beds, represented by the significantly disparate architectures of ATS-1/Trinity and ATS-2/Sierra. This will advance ASC co-design efforts by impacting one or more of the following:

- Determining performance bottlenecks in new architectures
- Characterizing ASC applications (for example, memory footprints, network utilization, instruction mixes) so as to better inform proxy application development and hardware design space evaluations
- Evaluating the portability and performance of different data structures and data layouts for important physics motifs across multiple architectures
- Evaluating the effectiveness of hardware and/or software-based memory and data management strategies to deal with automating movement between levels of system memory and I/O hierarchies
- Evaluating the performance analysis tools in these environments to determine their strengths and weaknesses, which will be invaluable in feedback to the vendors

Completion Criteria: This milestone will be completed when:

1. Improvements to the mini-applications to more closely represent our workloads are demonstrated on test-bed platforms and/or production systems.
2. A report has been completed by the 3 labs detailing lessons learned both successes and failures in regard to improvements.
3. The milestone team has communicated appropriate findings to vendors associated with the platforms evaluated in the milestone.

Impact Statement
The work performed under this milestone will lead to a single, more powerful, voice from the labs to the vendors when demanding changes in the tools and the toolchain and will have a
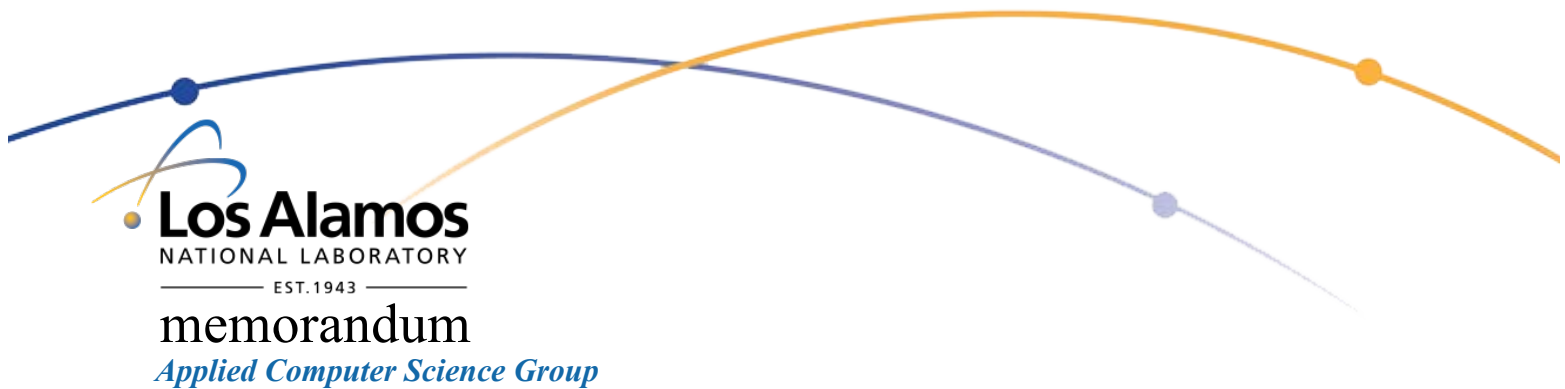
tangible impact on these relations moving forward. This is one of the first times that these tools have been applied to large applications and the results of these analyses will help improve both the applications themselves and the tools.

The work performed under this milestone should be a strong demonstration to vendors (and others) that while proxy-applications are a good starting point for studies, they are not enough and full applications need to be used as well since they oftentimes present new and tougher challenges to the tools.

One example of co-design with vendors that occurred during this milestone is the joint MPI work undertaken to alleviate increased resource pressure with larger number of ranks. The net result was a more usable MPI, which helps with all applications on the machine. This work demonstrated that the problem size is crucial when analyzing applications and that bottlenecks sometimes shift with problem sizes. This is something that should be conveyed to the vendors.

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September 11, 2017

From: Sriram Swaminarayan, LANL (Contact)
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To: Mike Lang, LANL
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Subject: 2017 Tri-lab ASC ATDM/CSSE/IC Co-design Level 2 Milestone

**Milestone Title: Tri-Lab Co-Design Milestone: In-Depth Performance Portability Analysis
        of Improved Integrated Codes on Advanced Architecture (#6016)**

Review Committee Comments:

The review committee believes the tri-lab co-design team fully met the completion criteria for
this milestone.

The work performed under this milestone will lead to a single, more powerful, voice from the
labs to the vendors when demanding changes in the tools and the toolchain and will have a
tangible impact on these relations moving forward.

The committee was impressed by the large body of work undertaken and the information
collected and presented in the slides. The committee was also impressed by the broad collection
of production codes that were used in this milestone – this is one of the first times that these tools
have been applied to large applications and the results of these analyses will help improve both
the applications themselves and the tools.

The work performed under this milestone should be a strong demonstration to vendors (and
others) that while proxy-applications are a good starting point for studies, they are not enough

and full applications need to be used as well since they oftentimes present new and tougher challenges to the tools.

One excellent example of co-design with vendors that occurred during this milestone is the joint MPI work undertaken to alleviate increased resource pressure with larger number of ranks. The net result was a more usable MPI, which helps with all applications on the machine.

This work demonstrated that the problem size is crucial when analyzing applications and that bottlenecks sometimes shift with problem sizes. This is something that should be conveyed to the vendors.

The committee found the discussion of worthiness of metrics useful and appropriate – if we cannot believe the output of the tools that we have, the hours (or days) spent collecting data are for naught. Some of the noteworthy examples presented included the interpretation of memory bandwidth reported in VTune and analysis of vector operations counts.

Finally, for the FY18 milestone, the committee would like to see more GPU results, stronger collaboration between the labs on metrics gathering, consistent metrics reporting across applications and labs, and inclusion of applications with memory accesses that have graph-like characteristics. We also recommend that the team share results with stakeholders, including experts that are knowledgeable in the relationship between the proxy application and the full application, early to ensure that proxy application results reflect expectations from the full application.

Sincerely,

Sriram Swaminarayan, on behalf of the committee.
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2017 Co-Design L2 Milestone Annotated Slides

Performance Tools and Ardra

David Richards
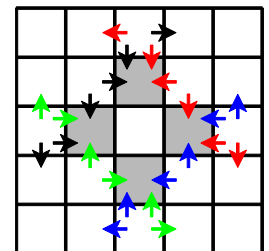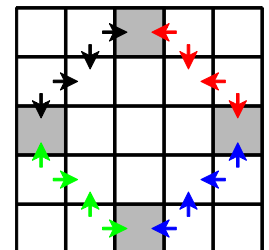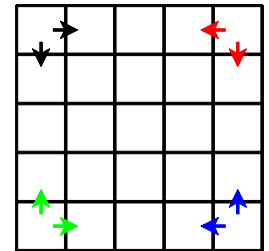ASC Co-Design Project Deputy

August 30, 2017

Lawrence Livermore National Laboratory

Ardra is LLNL's next generation deterministic transport code

- Solves the neutral particle Boltzmann Transport Equation in 1D, 2D, and 3D

- RAJA refactoring is in progress

- Solution is neutron energy, direction of flight, and spatial distribution
  - We have MPI parallelism in each of these dimensions

- 4 major kernels
  - LTimes, LPlusTimes, Sweep, Fission

- Jez_scale_3d
  - Critical bare Pu sphere
  - Weak scale with mesh cells (increase resolution)

- Optimization Challenges
  - Vectorization
  - Difficult memory access patterns
  - Hoisting index calculation
  - Peeling indices while descending call stack complicates exposing parallelism
  - Wavefront (sweep) dataflow

Kripke, a proxy app for Ardra, was used to explore use of RAJA

- Led to new RAJA N-nested loop abstraction

RAJA style nested for-loop

```
RAJA::View vview(v_ptr,
            make_perm_layout(ni,nj));
//...

RAJA::forallN< exec_policy, INDX, JNDX >(
            RangeSegment(1, ni),
            RangeSegment(0, nj),
            [=](INDX i, JNDX j) {

   vview(0, j) += vview(i, j);


});
```
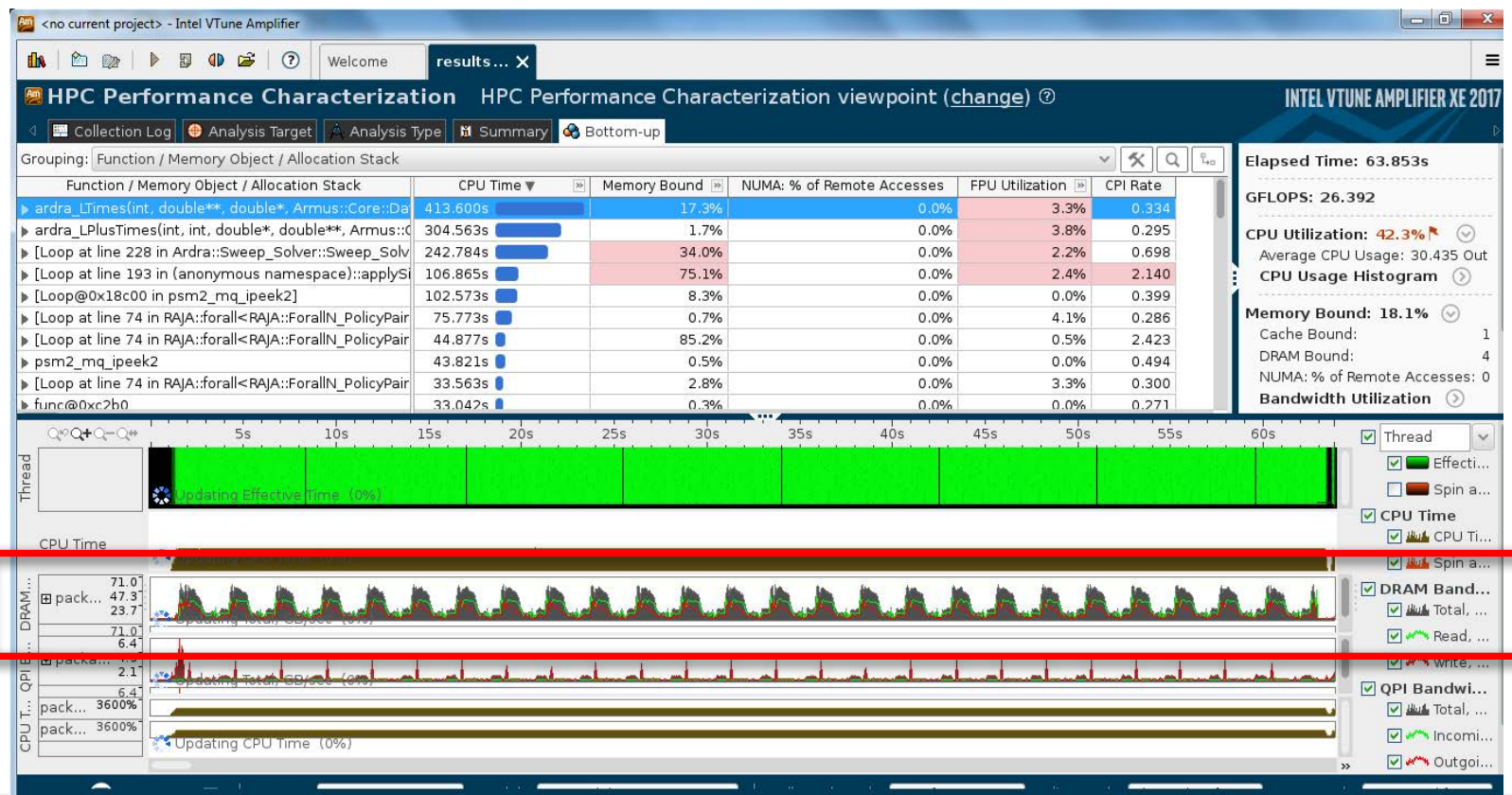
- Lessons learned from Ardra refactoring are driving new improvements to RAJA

- LTimes and LPlusTimes kernels are streaming kernels of the form phi(nm,g,z) += ell(nm,d) * psi(d,g,z)
  — z is stride 1
  — ell(nm,d) is invariant in inner loop

- Previous Kripke studies found LTimes & LPlusTimes (and all other kernels) are memory bound on GPU, but not elsewhere (less than 25% of B/W).

- *Goal: Identify the bottleneck for these kernels on current platforms*

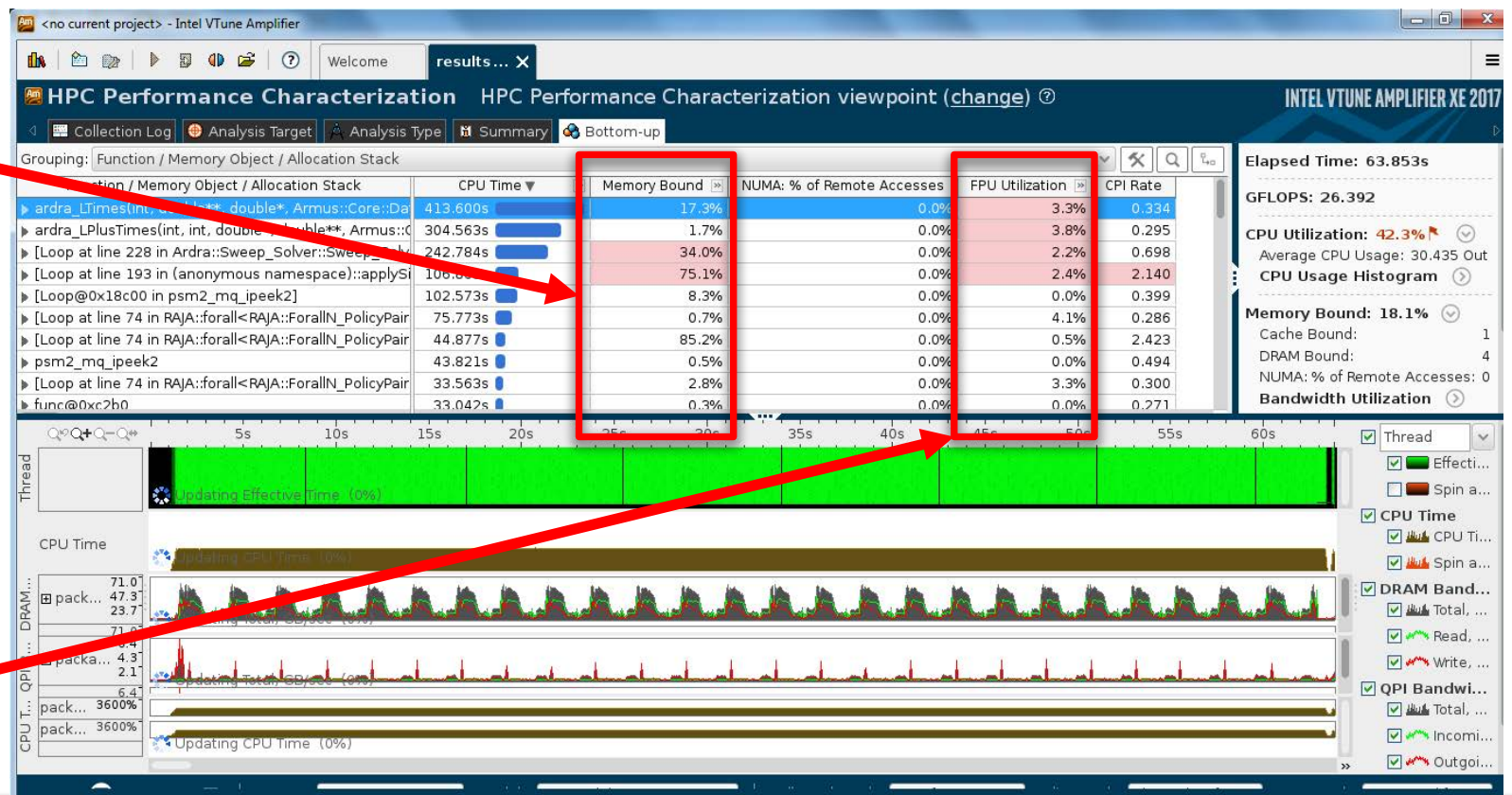# Intel VTune suggests Ardra is not B/W bound on CTS

DRAM B/W peaks at 50-60 GB/sec or ~50% of 130 GB/sec peak
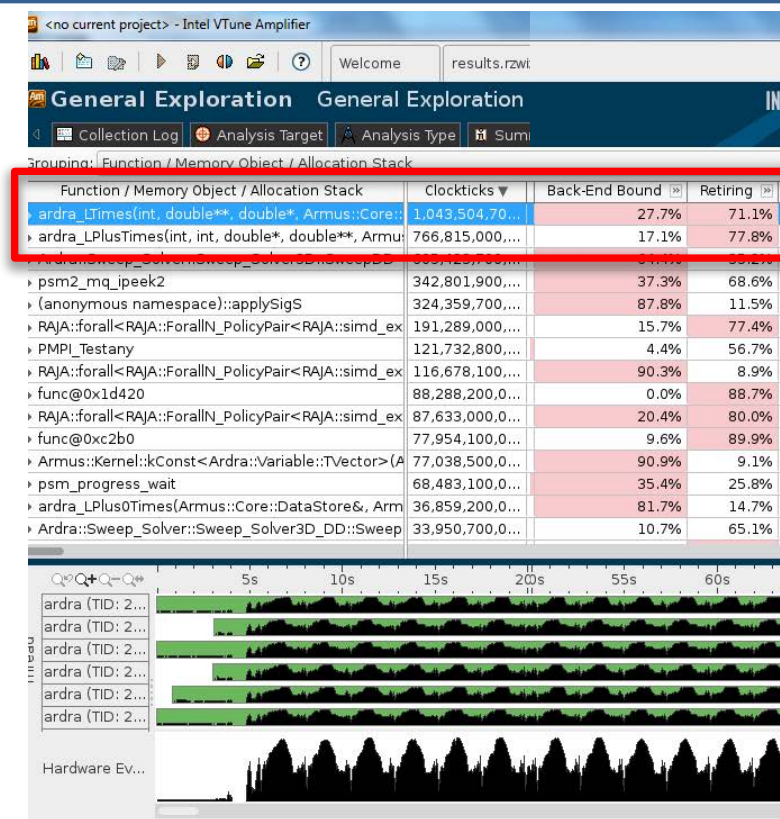
Intel VTune suggests Ardra is not B/W bound on CTS

This diagnostic is confusing and not consistent with developer intuition

FPU usage is uniformly low

Root cause analysis requires deduction and guesswork

- Low flop rate throughout code (2.1% of peak)

- Nearly all flops are scalar instructions
  - VTune reports 12:1 ratio for scalar to vector flops
  - Allinea Map reports 6:1 ratio
  - VTune shows no vectorization in key kernels

- VTune shows high "Retiring" in LTimes and LPlusTimes

- Intel optimization manual says high retirement means you should vectorize
  - Allows more operations to complete per instruction

After a lengthy saga, vectorization improved performance
(except when it didn't, and not the way we expected)

- RAJA abstractions are inhibiting vectorization
  - Refactoring and refinements are in progress
  - Manual unrolling and intrinsics were used to obtain vectorized code

- Performance improved:
  - Kernel times are significantly faster
  - But DRAM B/W changed very little

- Performance impact depends on problem size What's going on?
  - Small problem now operates at L2 B/W
  - Loads from L2 don't show up in DRAM B/W
  - Would blocking for L2 improve big problem performance?

Small Problem	scalar	vector
LTimes	15.53 sec	6.67 sec
LPlusTimes	22.43 sec	14.80 sec
DRAM B/W	60 GB/sec	68 GB/sec

Big Problem	scalar	vector
LTimes	84.75 sec	82.43 sec
LPlusTimes	101.38 sec	99.37 sec
DRAM B/W	117 GB/sec	119 GB/sec

CTS summary slide?

- Problem size matters

- B/W may not be what you think it is

- Bottleneck identification is hard

Ardra performance on KNL is ~1/2 of Broadwell (node to node)

- Limited tool set available. Mostly timing breakdowns by function
  - Data from internal timers
  - Similar data from CrayPAT, Map, PAPI

- KNL performance is consistent with expectations for codes not specifically tuned for KNL

- Performance difference not uniformly distributed across kernels

Time in Seconds	Broadwell	KNL	Ratio
Total	60.2	117.6	1.95
Solve	59.4	99.4	1.67
Sweep	12.3	31.8	2.58
LTimes	23.8	34.7	1.46
LPlusTimes	13.4	16.1	1.20
Fission	5.3	9.4	1.77

Ardra does not appear to be B/W bound on KNL

By changing NUMA modes we can get some hints that some kernels are B/W bound (at least they are sensitive), but we can't get at the details without better tools

Time in Seconds	W/O HBM	With HBM
Total	133	117.6
Solve	117.4	99.4
Sweep	31.0	31.8
LTimes	48.6	34.7
LPlusTimes	16.1	16.1
Fission	10.1	9.4

LTimes is sensitive, but not even close to a 4x speedup

And LPlusTimes shows no change?!

We hoped vectorization would improve KNL performance. Preliminary trials show only a 10% gain in LTimes.

Performance Co-Pilot helps provide more accurate memory bandwidth attribution

- Measuring memory B/W is challenging for multiple reasons
  - Counters are on memory controller ("uncore")
  - Not associated with any processor, thread, or execution context
  - Must use some model to attribute memory use to code

- Tools benefit from information from the application domain
  - PCP allows users to read memory controller counters and build their own models

- Data shown here assumes SPMD execution across all cores.

	VTune Memory Bound	PCP DRAM R/W bandwidth
Sweep	34%	20.0/12.5 GB/s
LTimes	17.3%	45.4/24.5 GB/s
LPlusTimes	1.75%	12.9/4.5 GB/s
Fission		47.1/45.9 GB/s

Ardra on GPU is relatively immature, but improving rapidly

- Kripke helped show RAJA would work in Ardra. Now Ardra is driving improvements in RAJA design and implementation

- RAJA N-nested loops re-designed/re-implemented to fix disastrous performance
  - CHAI copy constructors were being called for each iteration of the inner loop

- Trials to date obtain only ~5 GB/sec out of 732 GB/sec peak B/W
  - Clearly not memory bound. Larger problems are needed
  - Problem size is limited by RAJA reduction implementation
  - Fixes are in progress

- MPI communication in sweep kernel is a major bottleneck
  - New algorithm has been demonstrated
  - Not yet implemented in Ardra

Nvprof shows sweep is dominated by non-compute time

Manually scrolling and zooming is not the most productive method to identify such issues

Compare sweep to L0Times that shows dense compute

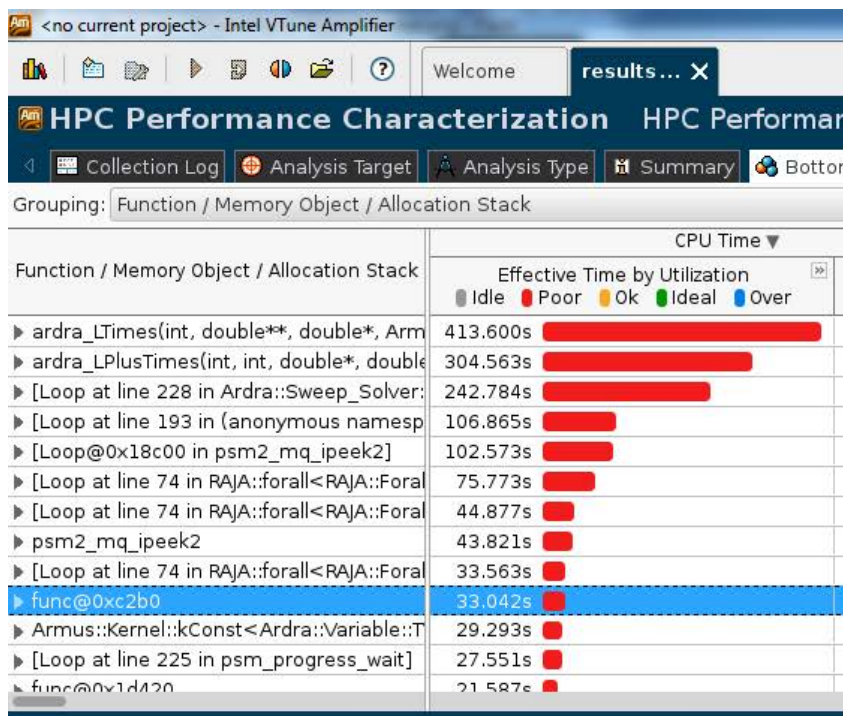# Other curiosities/frustrations: Tools don't agree on timings

Region	Region	Function
critSolver 97.5% 98.4% N/A 98% N/A		SigmaFTimes 10.7% 8.1% N/A N/A 1.0%
	Scattering 43.3% 55.6% 54.7% 64% 59.7%	LTimes 19.0% 23.9% 25.5% 40% 32.3%
		LPlusTimes 15.4% 20.0% 18.1% 23% 19.1%
		applySigS 7.5% 9.8% 11.1% N/A 8.3%
	Precond Solve 40.2% 30.7% N/A 35% N/A	MPI_Allreduce 20.9% 8.1% N/A N/A N/A
		Sweep_Solver3D 18.7% 21.9% 15.4% 20% 12.5%
Remainder 2.5% 1.6% N/A 2% N/A	Remainder 5.8% 5.6% N/A 1% N/A	Remainder 7.8% 9.2% N/A N/A N/A

Legend: HPCToolkit, MAP, GPerfTools, Built-In Timers, Open|SpeedShop
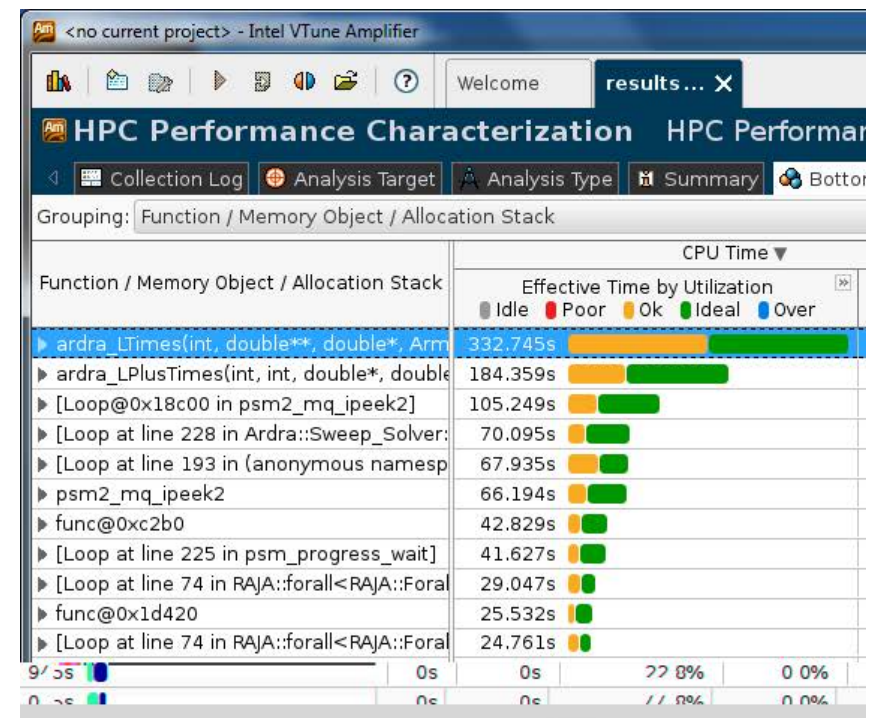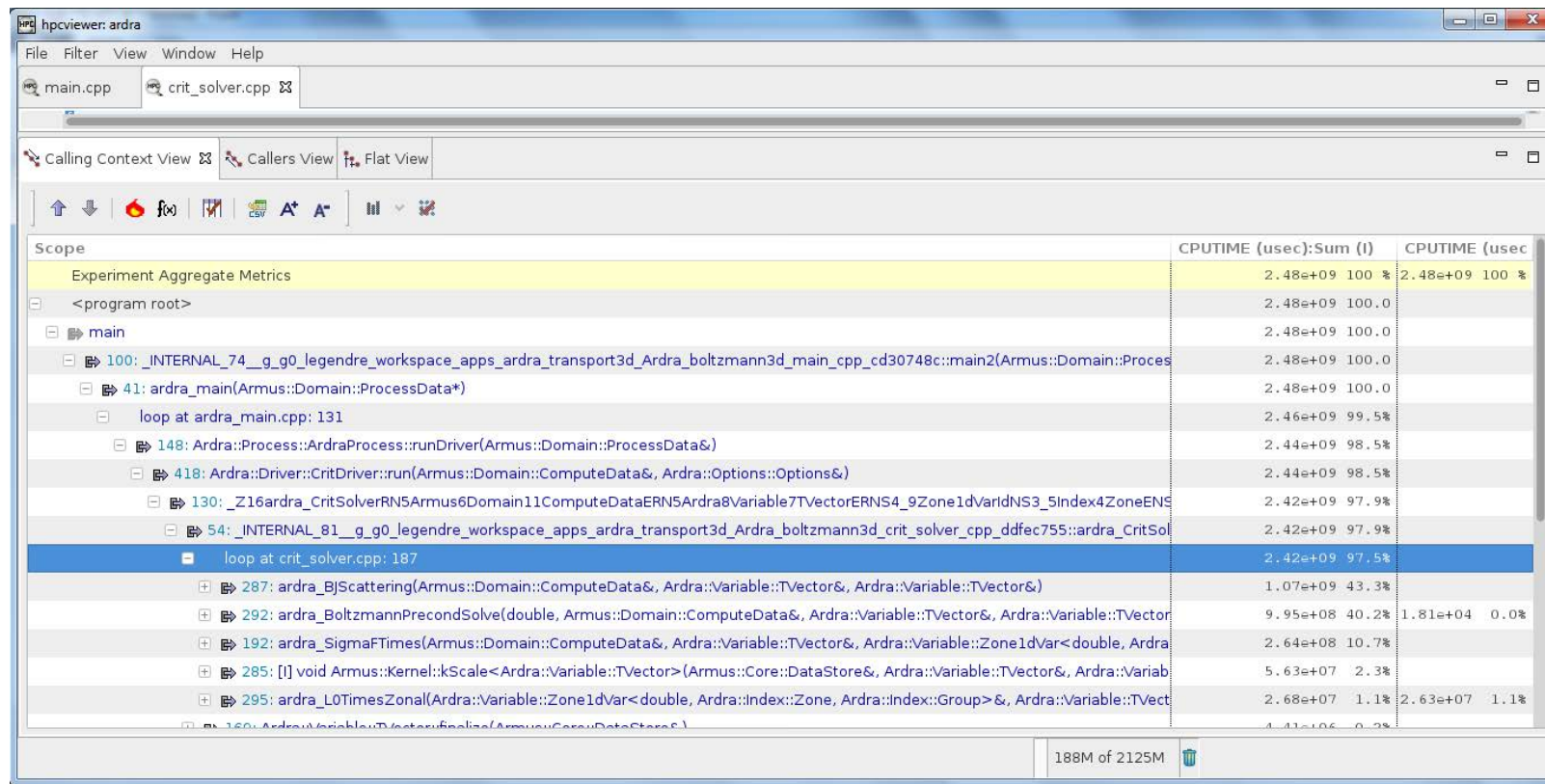
Other curiosities/frustrations: VTune recommendations can be less than helpful

32 threads: 55.75 sec

64 threads: 55.67 sec

Conclusion (needs to be refreshed/updated to latest results)

- CTS tools good enough for hypothesis generation
  - However, all guesses so far are wrong
  - Frequently end up looking at disassembled code

- Issues with misleading and confusing information
  - Incorrect or misleading hints are a huge waste of time

- Need to know algorithmic characteristics to fully understand performance
  - Especially memory reads/writes
  - Performance model and map to hardware are critical to understand if you're done optimizing

- Ardra is not B/W bound on any studied platform
  - GPU needs bigger problems
  - Bottleneck unclear on CTS and KNL

Backup slides showing that we have data for all of the metrics

Wall-Time Differences between Tools

Region	Region	Function
critSolver 97.5% 98.4% N/A 98% N/A		SigmaFTimes 10.7% 8.1% N/A N/A 1.0%
	Scattering 43.3% 55.6% 54.7% 64% 59.7%	LTimes 19.0% 23.9% 25.5% 40% 32.3%
		LPlusTimes 15.4% 20.0% 18.1% 23% 19.1%
		applySigS 7.5% 9.8% 11.1% N/A 8.3%
	Precond Solve 40.2% 30.7% N/A 35% N/A	MPI_Allreduce 20.9% 8.1% N/A N/A N/A
		Sweep_Solver3D 18.7% 21.9% 15.4% 20% 12.5%
Remainder 2.5% 1.6% N/A 2% N/A	Remainder 5.8% 5.6% N/A 1% N/A	Remainder 7.8% 9.2% N/A N/A N/A

Legend: HPCToolkit, MAP, GPerfTools, Built-In Timers, Open|SpeedShop

Ardra's Built-In Timers give an aggregated wall-time trace

```
-----------------------------------------------------------------------------
Run Time Statistics
                                                        Wall Time (sec)
%Tot %vl            Timer Name              Count      Total     Per Count
---- -------------------------------------- ----------- ---------- ----------

0.99   (0.99)  runDriver                          1  59.723425  59.723425
0.00     (0.00)  calcDetectorValues               2   0.000010   0.000005
0.00     (0.00)  fluxEdit                         1   0.000510   0.000510
0.99     (0.99)  solve                            1  59.401240  59.401240
0.99       (1.00)  critSolver                     1  59.401219  59.401219
0.38         (0.38)  ardra_BoltzmannPrecondSolve 25  22.586105   0.903444
0.20           (0.54)  Sweep3D_DD              1200  12.280318   0.010234
0.02           (0.02)  ardra_L0TimesZonal        25   0.947570   0.037903
0.63           (0.64)  bj_scattering_fanout      25  38.082054   1.523282
0.40             (0.63)  bj_gp_loop            1200  23.807309   0.019839
0.22             (0.35)  bj_lplus                25  13.394594   0.535784
0.00             (0.00)  dsaMatrixSetup           1   0.000970   0.000970
0.09           (0.09)  fission_fanout            26   5.359588   0.206138
0.00     (0.00)  writeRestart                     1   0.220500   0.220500
```

HPCToolkit can provide CPUTime, Wall-Time, PAPI Counters, MPI Time, Loop-level performance analysis
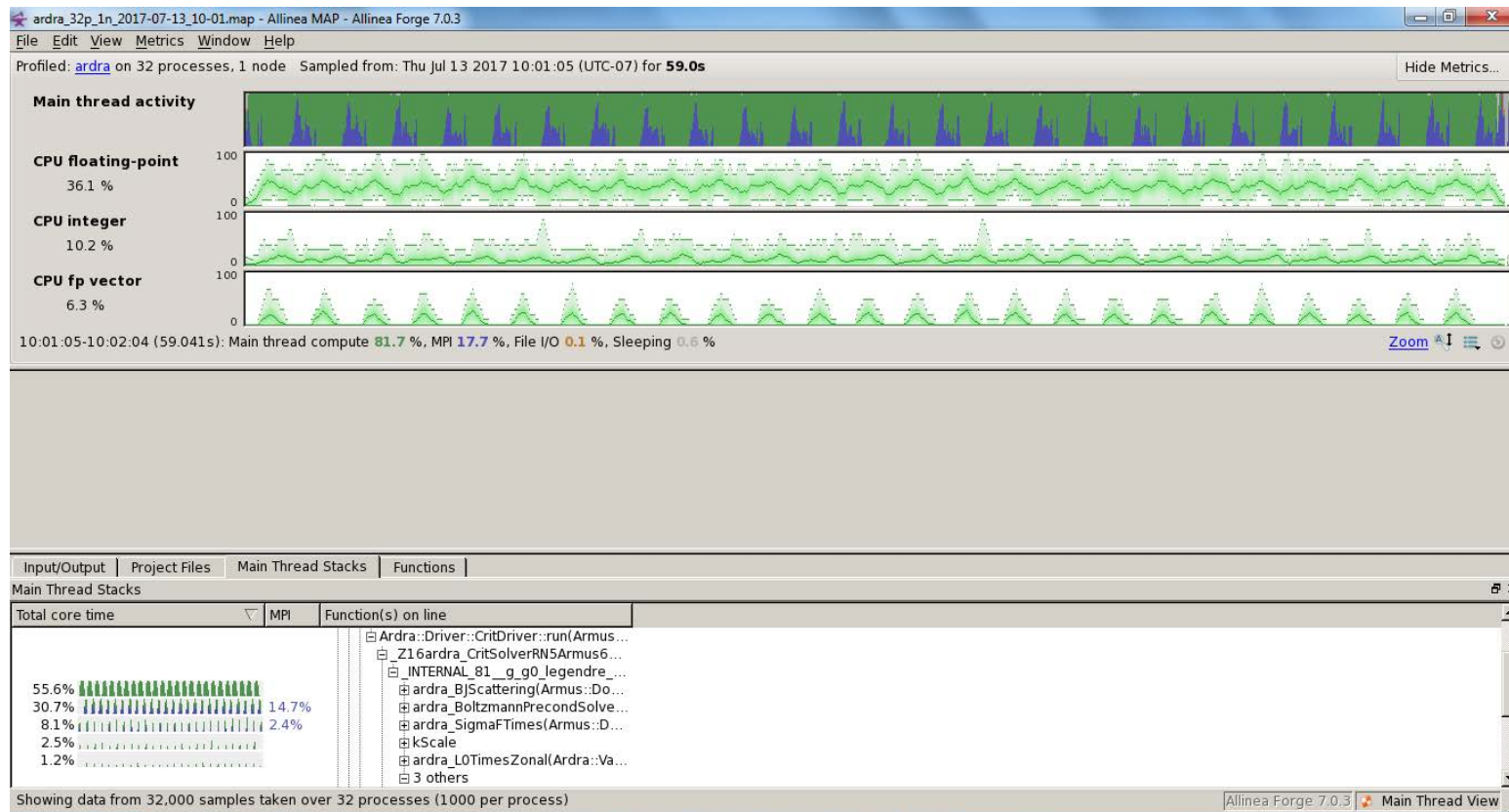
Google's GPerfTools provides an easy-to-use time breakdown

```
Total: 5536
1409  25.5%  25.5%      1410   25.5% ardra_LTimes::{lambda#2}::operator (inline)
 993  17.9%  43.4%      1000   18.1% ardra_LPlusTimes::{lambda#1}::operator (inline)
 855  15.4%  58.8%       855   15.4% Ardra::Sweep_Solver::Sweep_Solver3D::SweepDD
 483   8.7%  67.6%       613   11.1% ::applySigS
 477   8.6%  76.2%      3198   57.8% RAJA::forall (inline)
 324   5.9%  82.0%       324    5.9% Armus::Kernel::kGenericElemental2
 101   1.8%  83.9%       101    1.8% __intel_avx_rep_memset
  98   1.8%  85.6%       113    2.0% RAJA::LayoutBase_impl::operator (inline)
  85   1.5%  87.2%        85    1.5% PMPI_Testany
```

Allinea MAP Provides Wall Time, MPI Time, CPU Time, instruction mixes, FP vectorization rates
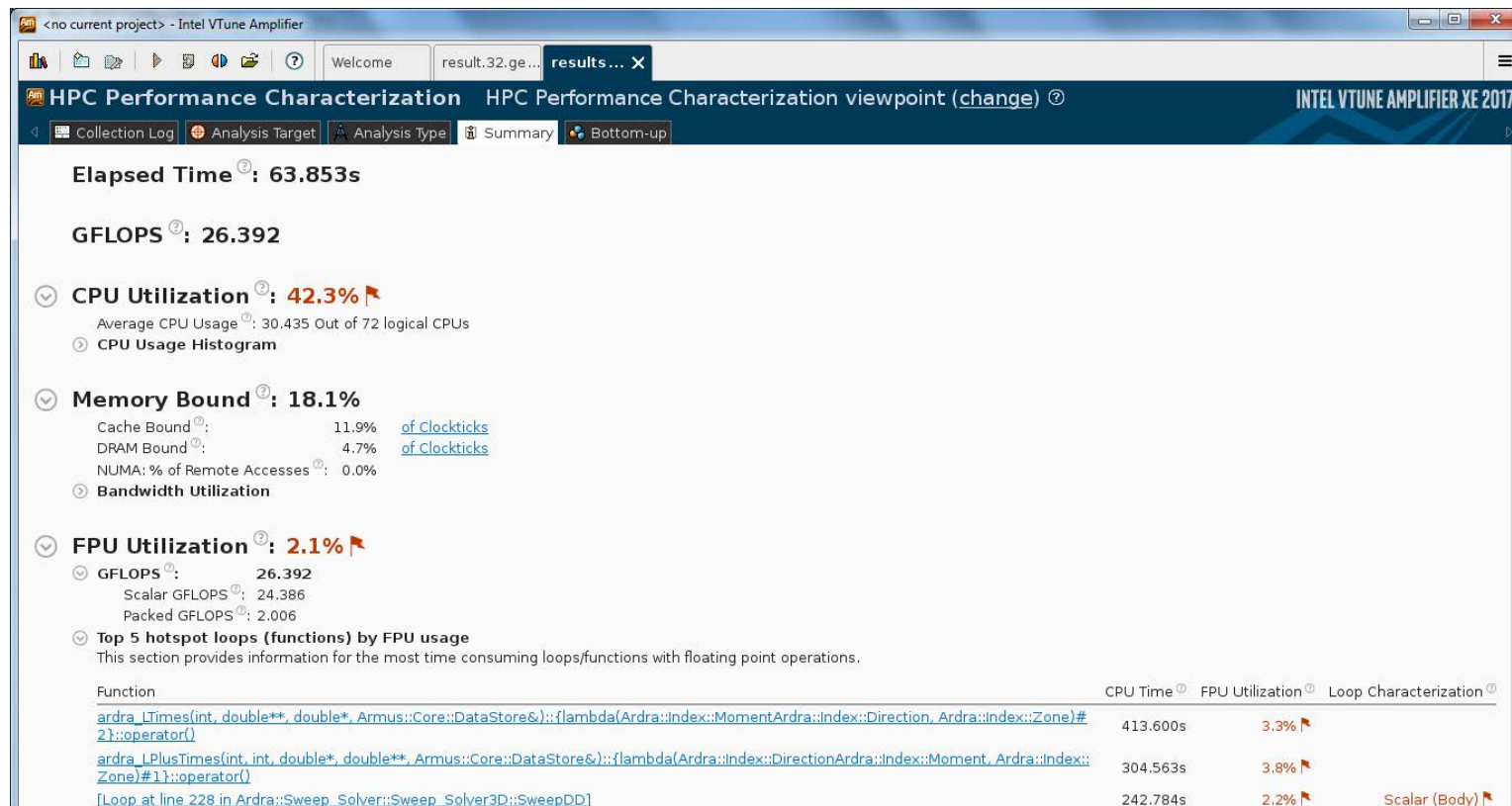
VTune and Vectorizing Ardra

- VTune Suggested we should vectorize the LTimes and LPlusTimes kernels (see next few slides)

- Vectorizing LTimes was a significant effort.

- Vectorization significantly speeds up Ardra for non memory-bound problem sizes.

Intel VTune Suggests FPU Utilization Issues across Ardra

Intel VTune Suggests FPU Utilization Issues in LTimes and LPlusTimes kernels

VTune Shows No Vectorization Happening in Ardra

VTune Shows Instruction Retirement Issues in LTimes and LPlusTimes

Intel Optimization Manual Says High Retirement Means You Should Vectorize

Intel 64 and IA-32 Architectures Optimization Reference Manual section B.1.7:

A high Retiring value for non-vectorized code may be a good hint to vectorize the code. Doing so essentially lets more operations to be [SIC] completed by single instruction/micro-op; hence improve performance.

VTune Vectorization Summary

- **Hypothesis**: We should vectorize LTimes and LPlusTimes:

- VTune Provided Evidence:
  — Low FPU Utilization in LTimes and LPlusTimes
  — No Vector instructions being compiled into LTimes and LPlusTimes
    - Also backed by MAP and disassembling the kernels
  — LTimes and LPlusTimes spends most time retiring instructions
  — Intel optimization manual says retiring is optimized via vectorization

- Used Intel compiler intrinsics to vectorize LTimes into a loop around a 256-bit fused add/multiply instruction.

- **Results:** Vectorizing LTimes sped up by over 2x. Problem now runs at the speed of L2 cache

MPI Weak Scaling with Ardra

- The sweep algorithms used by Ardra are known not to scale

- Ardra has previously been shown to scale as well as the ideal sweep algorithm

- The mpiP tool can break down MPI times by call path and point-to-point vs. collective
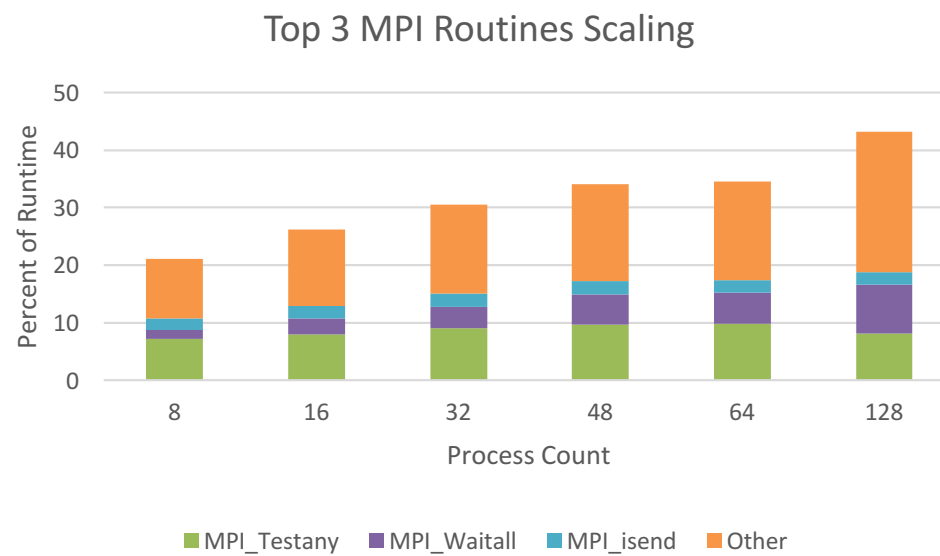
Scaling Top 3 MPI Routines in Ardra

Top 3 MPI Routines Scaling

Other Measurements

- Memory High Water Mark

- Memory Bandwidth

- IPC

memP Tool Shows Memory High-Water Mark

```
-------------------------------------------------------------------------
@--- Greatest Heap High Water Mark (top 5, descending, bytes)
-------------------------------------------------------------------------
        Rank        Heap HWM        Stack            Sum
         16        334.45 MB     27.55 KB      334.48 MB
          1        334.27 MB     27.64 KB      334.30 MB
          4        334.27 MB     27.64 KB      334.30 MB
         17        334.19 MB     27.55 KB      334.21 MB
         20        334.19 MB     27.55 KB      334.21 MB
-------------------------------------------------------------------------
@--- Heap HWM Statistics ------------------------------------------------
-------------------------------------------------------------------------
Max                               :     334.45 MB
Median                            :     310.12 MB
Mean                              :     310.61 MB
Min                               :     287.77 MB
Stddev                            :      16.32 MB
Coefficient of variation          :      0.052555
```

Performance Co-Pilot Provides Access to Memory Bandwidth Counters

```
Run Time Statistics                                Wall Time (sec)
     Name                          Count   Total     Per Count   Read BW (MB/S)   Write BW (MB/S)
---- ---------------------------- -------- --------- ---------   ----------------  --------------

0.99   (0.99)  runDriver              1   58.194303  58.194303        33521.11         19988.36
0.00    (0.00)  calcDetectorValues    2    0.001631   0.000816        58550.63          5699.95
0.00    (0.00)  fluxEdit              1    0.001138   0.001138        19093.88           449.64
0.98    (0.99)  solve                 1   57.421936  57.421936        33789.96         20242.19
0.98     (1.00)  critSolver           1   57.421919  57.421919        33787.91         20241.49
0.38      (0.39)  ardra_BoltzmannPrecondSolve  25  22.263067  0.890523   11457.50       6998.44
0.20       (0.53)  Sweep3D_DD        1200  11.754593   0.009795        20016.79         12478.46
0.01       (0.01)  ardra_L0TimesZonal  25   0.839731   0.033589       107391.46          1972.86
0.65      (0.66)  bj_scattering_fanout  25  38.104032  1.524161        33744.05         18284.05
0.39       (0.60)  bj_gp_loop        1200  23.008618   0.019174        45400.02         24469.75
0.22       (0.34)  bj_lplus            25  12.917722   0.516709        12885.62          4581.34
0.00      (0.00)  dsaMatrixSetup        1    0.000816   0.000816         8204.87          5446.62
0.08      (0.08)  fission_fanout       26    4.458054   0.171464        47140.92         45936.05
0.01    (0.01)  writeRestart           1    0.668223   0.668223          412.27           341.27


----------------------------------------------------------------------------
```

IPC calculations from PAPI

```
                                                 Wall Time (sec)
%Tot    %lvl    Timer Name              Count     Total   Per Count    PAPI_TOT_INS    PAPI_TOT_CYC    IPC
----  --------  --------------------------- -----------  ----------  ------------  --------------  ----
0.99  (0.99)  runDriver                     1  55.470365  55.470365   324162081674    137489604904   2.36
0.00    (0.00)  calcDetectorValues          2   0.000031   0.000015           4766           11925   0.40
0.00    (0.00)  fluxEdit                    1   0.000435   0.000435        1416821         1026943   1.38
0.98    (0.99)  solve                       1  54.835454  54.835454   323421631027    137001207761   2.36
0.98      (1.00)  critSolver                1  54.835402  54.835402   323421624234    137001189017   2.36
0.35        (0.35)  ardra_BoltzmannPrecondSolve 25 19.416221 0.776649 462818928       28673595986   1.61
0.20          (0.58)  Sweep3D_DD         1200  11.347921   0.009457    41570458833     26857361098   1.55
0.01          (0.01)  ardra_L0TimesZonal 25   0.817820   0.032713     6106458767      2090150611   2.92
0.64        (0.66)  bj_scattering_fanout 25  35.924563  1.436983    245742185602     90729230958   2.71
0.40          (0.62)  bj_gp_loop       1200  22.161912   0.018468   133669003312     56761088769   2.35
0.23          (0.35)  bj_lplus           25  12.657956   0.506318   111588764852     32601102609   3.42
0.00        (0.00)  dsaMatrixSetup        1   0.000040   0.000040           2511            9463   0.27
0.08        (0.08)  fission_fanout       26   4.409059   0.169579    22592708167     10981676930   2.06
0.01      (0.01)  writeRestart            1   0.534155   0.534155      358056618       236388890   1.51
```
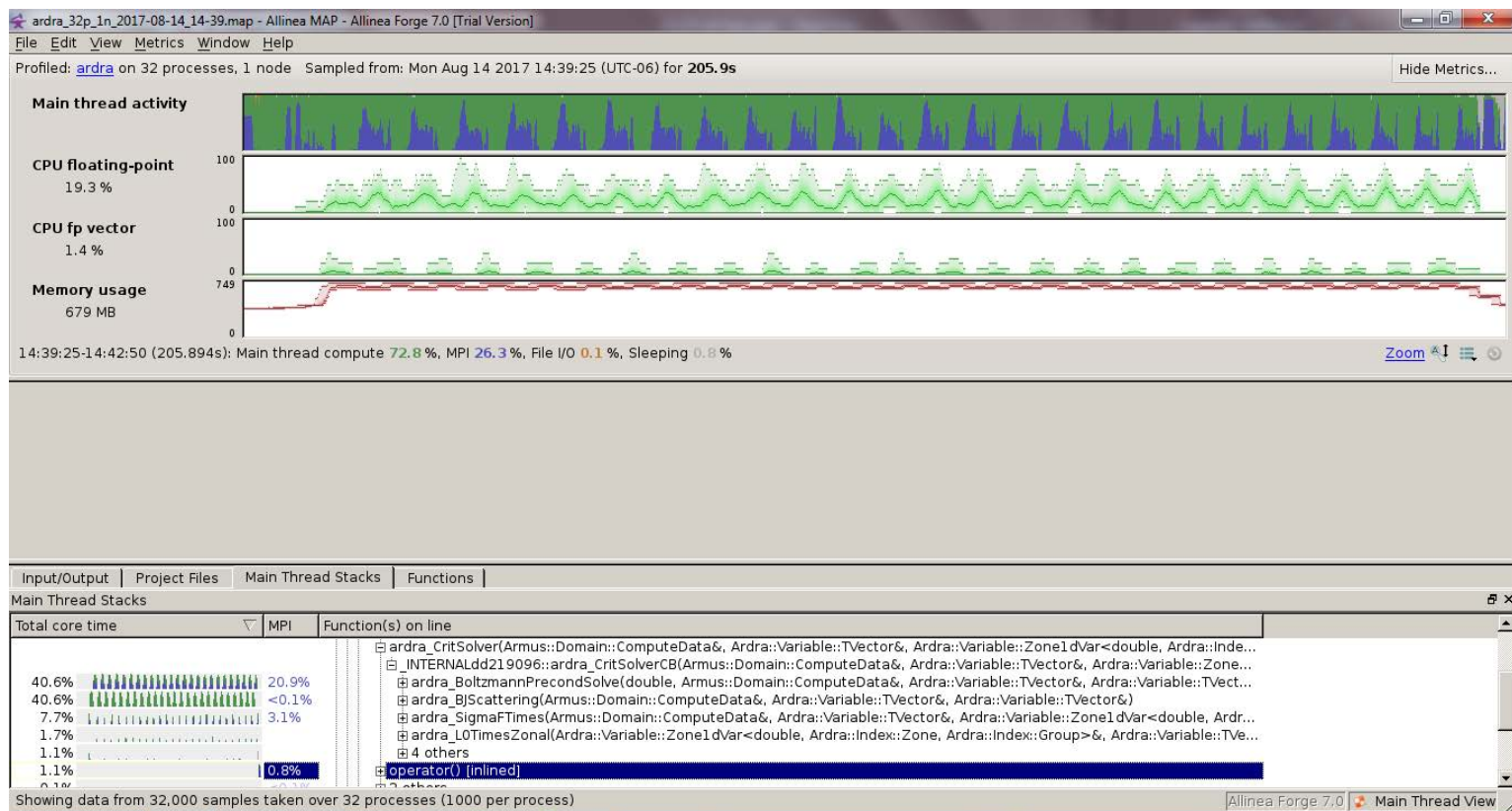
KNL Data for Ardra

- Run on LANL's Trinitite

Allinea MAP on KNL Showing Instruction Mix

KNL CrayPAT Hardware Counters

```
Table 4:  Program HW Performance Counter Data

PE=HIDE


============================================================================
   Total
----------------------------------------------------------------------------
   UNHALTED_CORE_CYCLES             148,513,518,999
   UNHALTED_REFERENCE_CYCLES       138,612,610,990
   INSTRUCTION_RETIRED             121,483,923,581
   LLC_REFERENCES                    2,923,544,676
   LLC_MISSES                           91,517,871
   LLC cache hit,miss ratio  96.9% hits        3.1% misses
============================================================================
```

Ardra's Built-In Timers, with HBM

```
--------------------------------------------------------------------------------
Run Time Statistics
                                                            Wall Time (sec)
%Tot %vl          Timer Name              Count      Total      Per Count
---- ---------------------------------------- ----------- ---------- ----------
0.86  (0.86)  runDriver                            1 101.049208 101.049208
0.00    (0.00)  calcDetectorValues                 2   0.000017   0.000009
0.00    (0.00)  fluxEdit                           1   0.014949   0.014949
0.84    (0.98)  solve                              1  99.364175  99.364175
0.84     (1.00)  critSolver                        1  99.364004  99.364004
0.45      (0.54)  ardra_BoltzmannPrecondSolve     25  53.269412   2.130776
0.27       (0.60)  Sweep3D_DD                    1200  31.825563   0.026521
0.02       (0.02)  ardra_L0TimesZonal              25   2.202952   0.088118
0.44      (0.52)  bj_scattering_fanout            25  51.718966   2.068759
0.30       (0.67)  bj_gp_loop                    1200  34.733989   0.028945
0.14       (0.31)  bj_lplus                        25  16.131730   0.645269
0.00      (0.00)  dsaMatrixSetup                   1   0.000052   0.000052
0.08      (0.09)  fission_fanout                  26   9.390708   0.361181
0.01    (0.02)  writeRestart                       1   1.615085   1.615085
```

Ardra's Built-In Timers, with HBM disabled

```
--------------------------------------------------------------------------
Run Time Statistics
                                                        Wall Time (sec)
%Tot %vl           Timer Name            Count     Total      Per Count
---- ------------------------------------------- ----------- --------- ----------
0.88  (0.88)  runDriver                           1 117.428259 117.428259
0.00    (0.00)  calcDetectorValues                2   0.000019   0.000010
0.00    (0.00)  fluxEdit                          1   0.012143   0.012143
0.87    (0.99)  solve                             1 115.775249 115.775249
0.87      (1.00)  critSolver                      1 115.774777 115.774777
0.39        (0.45)  ardra_BoltzmannPrecondSolve  25  52.312439   2.092498
0.23          (0.59)  Sweep3D_DD               1200  31.040346   0.025867
0.02          (0.02)  ardra_L0TimesZonal         25   2.198163   0.087927
0.50        (0.57)  bj_scattering_fanout         25  66.402436   2.656097
0.36          (0.73)  bj_gp_loop               1200  48.644699   0.040537
0.12          (0.24)  bj_lplus                   25  16.083930   0.643357
0.00        (0.00)  dsaMatrixSetup                1   0.000026   0.000026
0.08        (0.09)  fission_fanout              26  10.131977   0.389691
0.01    (0.01)  writeRestart                      1   1.543792   1.543792
```

Ardra's Built-In Timers with IPC from PAPI

```
--------------------------------------------------------------------------------
Run Time Statistics

                                         Wall Time (sec)
%Tot %vl        Timer Name          Count    Total   Per Count   PAPI_TOT_INS    PAPI_TOT_CYC      IPC
---- ------------------------------ ----- ---------- ---------  --------------  --------------    ------
0.85  (0.85)  runDriver                 1  98.596496  98.596496  127939335250   136701452657      0.94
0.00    (0.00)  calcDetectorValues      2   0.000068   0.000034          4394          26922      0.16
0.00    (0.00)  fluxEdit                1   0.011763   0.011763        764245        1192409      0.64
0.84  (0.98)  solve                     1  97.051414  97.051414  127129798718   134946610048      0.94
0.84    (1.00)  critSolver              1  97.051194  97.051194  127129792939   134946537515      0.94
0.44      (0.52)  ardra_BoltzmannPrecondSolve  25  50.812436   2.032497   34422457161    55545714728      0.62
0.27        (0.62)  Sweep3D_DD        1200  31.608640   0.026341   29440200696    46646236415      0.63
0.02        (0.02)  ardra_L0TimesZonal  25   2.175385   0.087015    3168726418     3177950528      1.00
0.43      (0.52)  bj_scattering_fanout  25  50.008220   2.000329   77215472964    62496507175      1.24
0.29        (0.67)  bj_gp_loop       1200  33.539460   0.027950   46714375423    41527300851      1.12
0.13        (0.31)  bj_lplus           25  15.503903   0.620156   30440374680    20618413930      1.48
0.00      (0.00)  dsaMatrixSetup        1   0.000083   0.000083          2567          20830      0.12
0.08      (0.09)  fission_fanout       26   9.211234   0.354278   10944412869    12622144189      0.87
0.01  (0.02)  writeRestart              1   1.536780   1.536780     617062571     1529302924      0.40
```

Using the LTimes Vector Optimizations from CTS-1 has smaller impact on KNL

- LTimes when scalar: 31.68s

- LTimes after vector optimizations: 34.88s

GPU Raw Data

- These benchmarks were run on rzmanta, a Coral EA system at LLNL. Some key specs
  - Each node has 4 Tesla P100-SXM2 Nvidia GPUs
  - Also Power8 cores
  - Power8+GPU linked with NVLINK to enable
  - Final numbers gotten from CUDA 8.0, system now has CUDA9, numbers appear similar

GPU Bandwidth Measurements

Compute-Dense L0TimesZonal shows good utilization of P100s

Compute-sparse Sweep Shows poor utilization of P100s

Parallel utilization is low across the application, kernels are small

Utilization Numbers

IPC/CPI for GPU

Cache hit rates are relatively good on this problem size

Los Alamos

NATIONAL LABORATORY

— EST.1943 —

Delivering science and technology
to protect our nation
and promote world stability

Co-Design L2 Milestone Review

Ben Bergen

August 30th, 2017

- **Terminology**
- **PARTISN**
- **VPIC**
- **FleCSALE**

Pipeline Hazards

- **Structural Hazard**

  - A hazard that arises from a resource conflict where the hardware cannot support all possible instruction combinations simultaneously

  - Example: the pipeline stalls because a vector addition requires both of its operands to be shuffled, but shuffles can only occur on one execution unit
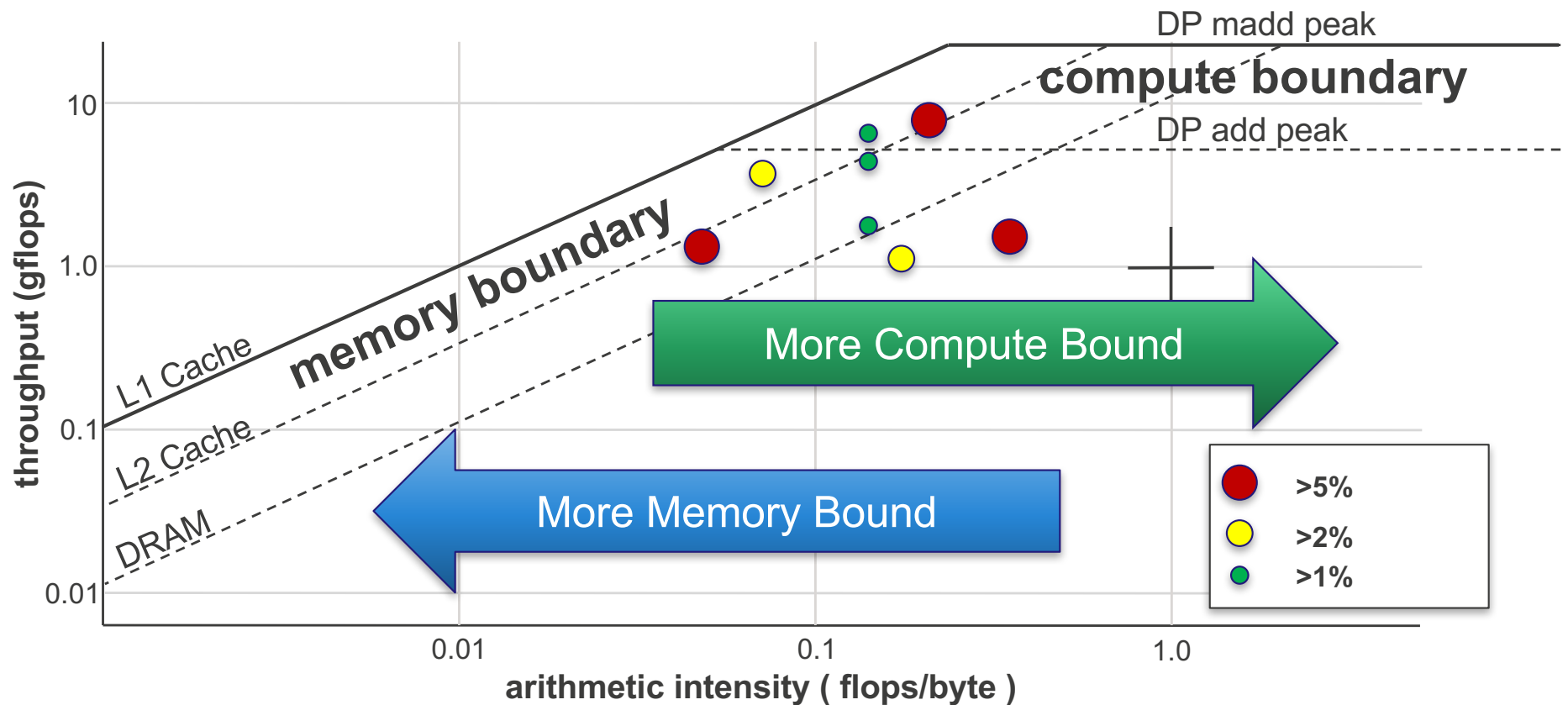
- **Data Hazard**

  - A hazard that arises when an instruction depends on the result of a previous instruction

  - Example: the pipeline stalls because a load operation from main memory has not yet completed

- **Control Hazard**

  - A hazard that arises from a branch or instruction that changes the program counter

Intel Roofline Model – Explanation

PARTISN
Work done by Randy Baker and Joe Zerr
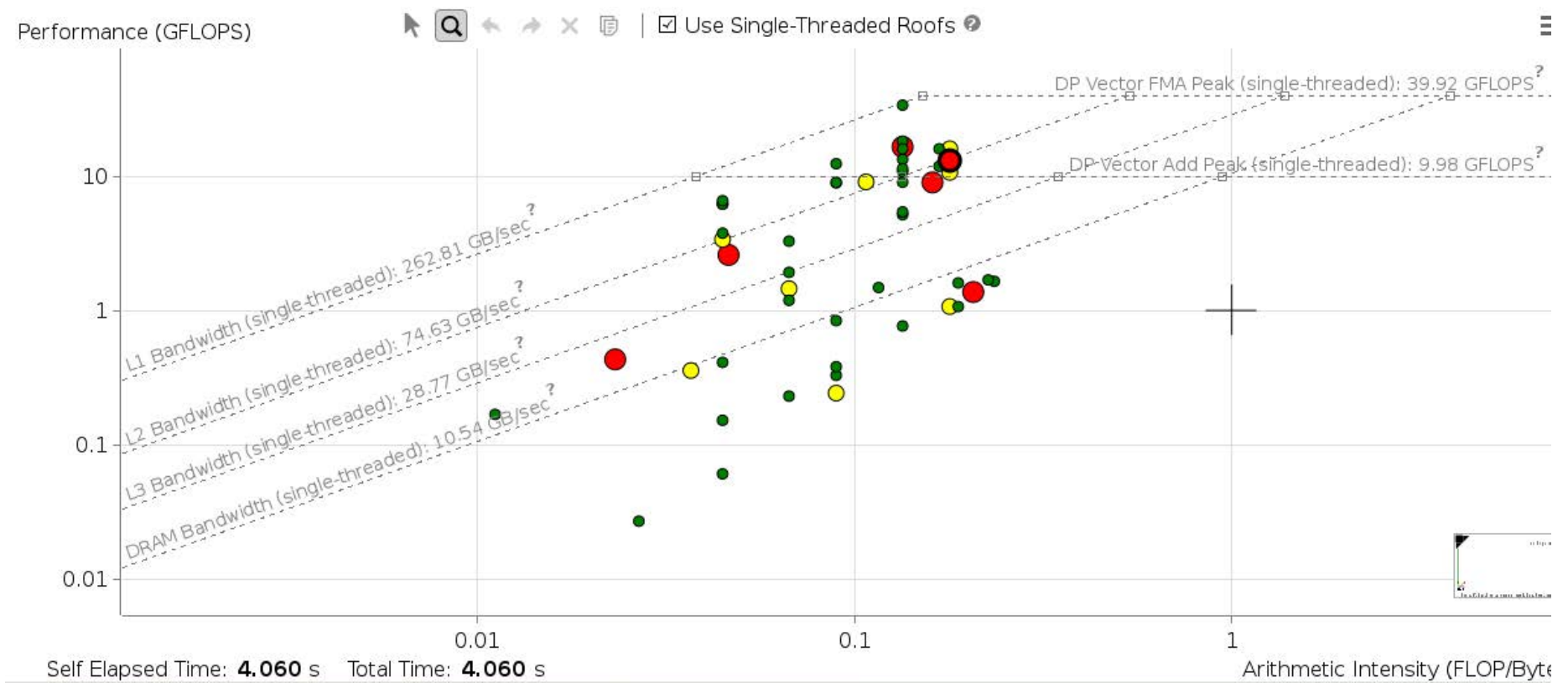
PARTISN
Overview

PARTISN is a time-dependent, multi-dimensional neutron/gamma transport code developed at LANL

- **Basis for SNAP mini-app**
- **110,000 lines of code (Fortran with some C and Python)**
- **Dates to the 1950s/early 1960s**
- **Originally targeted to the CDC-7600**
  - 10 MFlop/s peak
- **Supported by four people (2.5 FTEs)**
  - Methods research
  - Algorithm development
  - Software development & testing
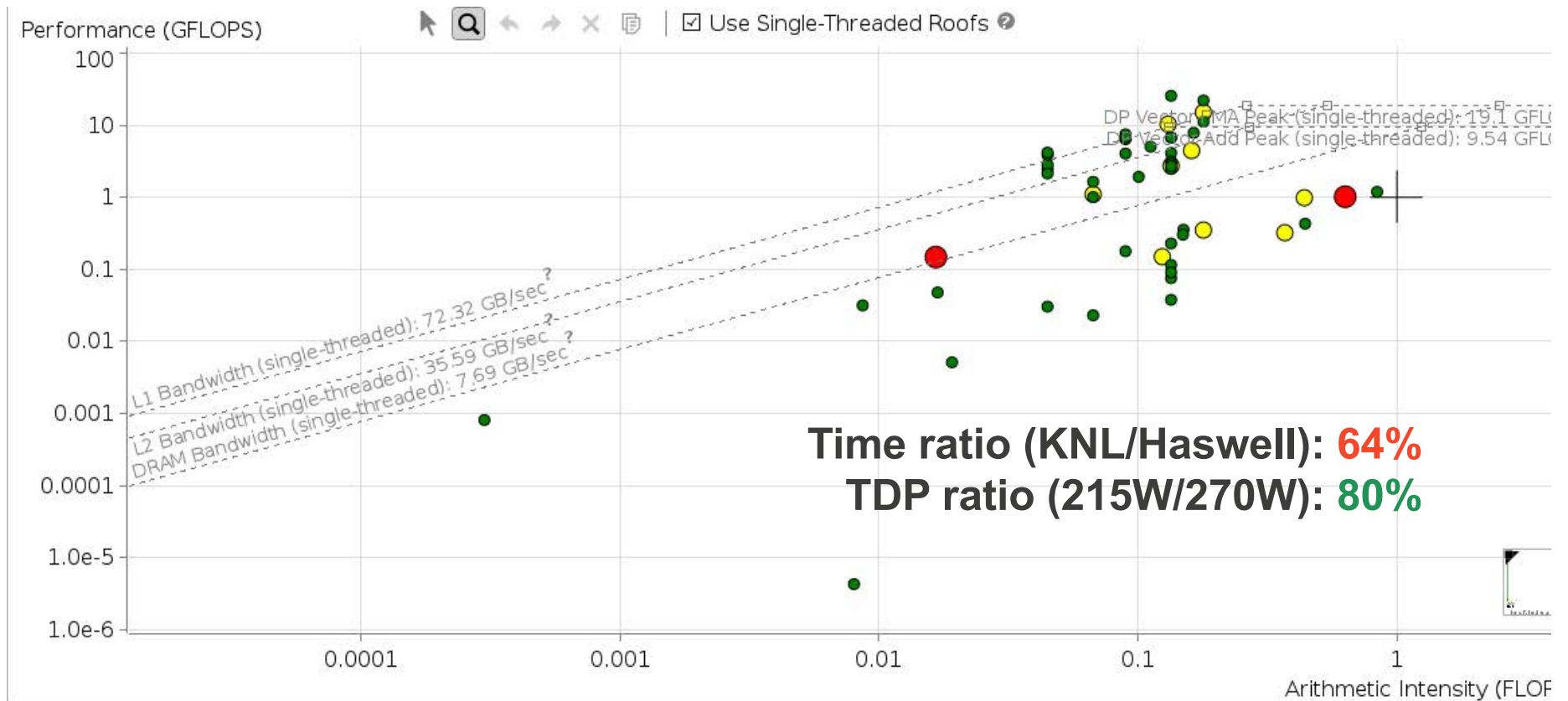 - User support & applications expertise

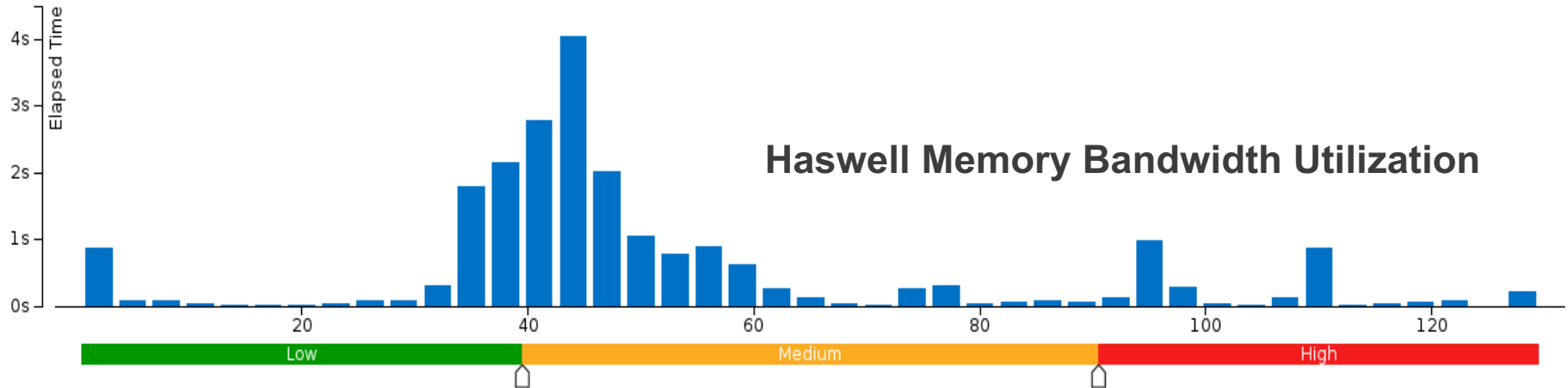# PARTISN
Haswell 32-core, 32 MPI ranks, 32 GB, 48.0 seconds

PARTISN
KNL 68-core, 64 MPI ranks, 32 GB, 74.6 seconds

Time ratio (KNL/Haswell): **64%**
TDP ratio (215W/270W): **80%**

PARTISN
Memory Bandwidth Boundedness

Haswell Memory Bandwidth Utilization

- **PARTISN is only *somewhat* memory bandwidth bound on Haswell**

  - 13.2% DRAM bandwidth bound

  - 0.519 CPI (minimum is 0.25 CPI)

- **Memory bandwidth is less of an issue on KNL**

  - 4.0% DRAM bandwidth bound

  - 1.326 CPI (minimum is 0.5 CPI)

PARTISN
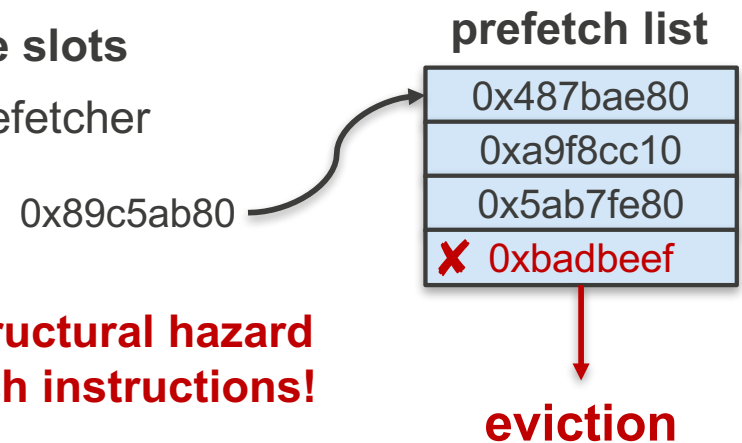# KNL Bottleneck – Prefetch

- **Intel VTune Amplifier**

CPU Time ⑦:	75.040s
▽ **Memory Bound:**	
L2 Hit Rate ⑦:	91.0%
L2 Hit Bound ⑦:	21.5% of Clockticks
L2 Miss Bound ⑦:	28.8% ⚑ of Clockticks
MCDRAM Cache Bandwidth Bound ⑦:	0.1%
MCDRAM Flat Bandwidth Bound ⑦:	0.3%
DRAM Bandwidth Bound ⑦:	3.6% of Elapsed Time
L2 Miss Count ⑦:	141,004,230
MCDRAM Hit Rate:	74.3%
MCDRAM HitM Rate:	46.5%

- **Red flag is L2 Miss Bound**
  - In spite of > 90% hit rate
- **Discussion with Intel suggested that compiler-generated prefetch instructions were swamping the prefetcher**

PARTISN
KNL Bottleneck – Prefetch

- **Prefetch instructions can reduce memory latency by requesting data before they are needed, giving more leeway to the dynamic scheduler**

- **KNL supports software and hardware-generated prefetch instructions (priority is given to software-generated prefetch instructions)**

- **The prefetcher has a limited number of available slots**
  - Too many prefetch instructions can *thrash* the prefetcher

prefetch list

0x89c5ab80

| 0x487bae80 |
| 0xa9f8cc10 |
| 0x5ab7fe80 |
| ✘ 0xbadbeef |

eviction

- **The likely cause of the KNL bottleneck was a structural hazard caused by too many compiler-generated prefetch instructions!**

PARTISN
KNL Optimization – Manual Prefetch

- **Disabled top-level loop prefetch**
  - *cdir$ noprefetch*
- **Added manual prefetch instructions for key arrays ahead of use**
  - *call mm_prefetch (q(n,i,j,k), 1)*
  - L2 prefetch is preferable to L1
  - data alignment is not critical
- **Labeled key array as non-temporal, i.e., don't bother caching this…**
  - *cdir$ vector aligned nontemporal*
  - data alignment is critical
- **Single day of actual coding effort**

PARTISN
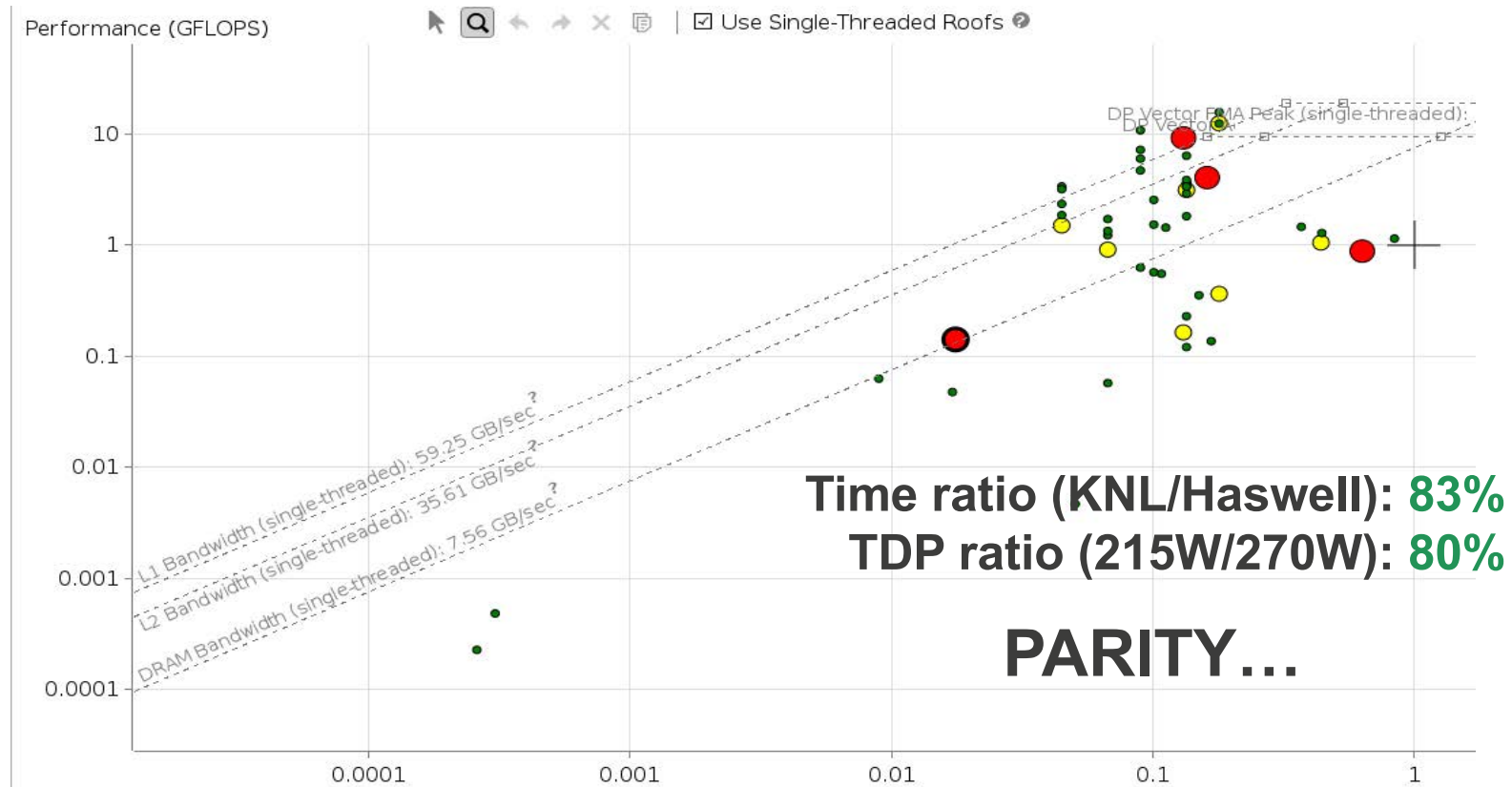# Optimizations Significantly Improved KNL Performance

- **Intel VTune Amplifier**

CPU Time ⓘ:	75.040s - 67.720s = 7.320s	
⊙ **Memory Bound:**		
L2 Hit Rate ⓘ:	91.0% - 96.0% = -5.0%	
L2 Hit Bound ⓘ:	21.5% - 34.8% = -13.3%	of Clockticks
L2 Miss Bound ⓘ:	28.8% - 19.7% = 9.1%	of Clockticks
MCDRAM Cache Bandwidth Bound ⓘ:	Not changed, 0.1%	
MCDRAM Flat Bandwidth Bound ⓘ:	0.3% - 0.4% = -0.2%	
DRAM Bandwidth Bound ⓘ:	3.6% - 4.0% = -0.4%	of Elapsed Time
MCDRAM Hit Rate:	74.3% - 74.6% = -0.2%	
MCDRAM HitM Rate:	46.5% - 47.4% = -0.9%	
Total Thread Count:	1 - 2 = -1	
Paused Time ⓘ:	Not changed, 0s	
L2 Miss Count ⓘ:	141,004,230 - 87,002,610 = 54,001,620	

- **L2 Hit Rate**
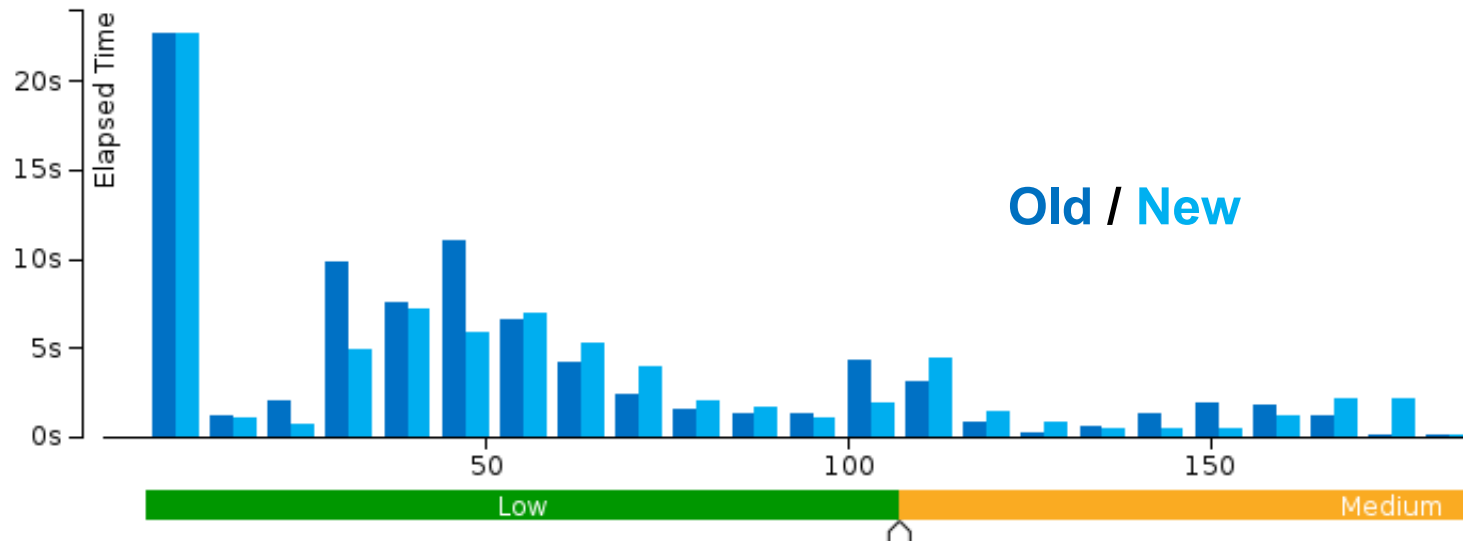 - Improved by 5%
- **L2 Miss Bound**
  - Decreased by 9%

PARTISN
KNL 68-core, 64 MPI ranks, 32 GB, 56.9 seconds

Time ratio (KNL/Haswell): **83%**
TDP ratio (215W/270W): **80%**

PARITY…

PARTISN
Memory Bandwidth Boundedness

Old / **New**

- **PARTISN is *still* only somewhat memory bandwidth bound**
- **The implication is that we still haven't achieved optimal performance…**
  - We are exploring scheduling issues for nested vector operations

PARTISN
Summary

- **PARTISN is not primarily memory bandwidth bound**
  - KNL performance optimization removed a structural hazard that was affecting memory latency
  - Remaining performance gains are likely to be gotten by improving scheduling of vector operations
- **Coral-EA**
  - Code has been ported to compile on LANL's pre-Sierra nodes (Power8)
  - Future work will investigate OpenMP function off-load capability

VPIC
Work done by Dave Nystrom

VPIC
Overview

VPIC is a 3D explicit, relativistic, charge-conserving electromagnetic particle-in-cell (PIC) code originally developed by Kevin Bowers at LANL

- **Open Source: https://github.com/losalamos/vpic**
- **Single Precision**
- **Structured Cartesian Mesh**
- **Array-of-Structures (AoS) (particle and field data)**
- **Asynchronous MPI (distributed memory)**
- **OpenMP or Pthreads (shared memory)**
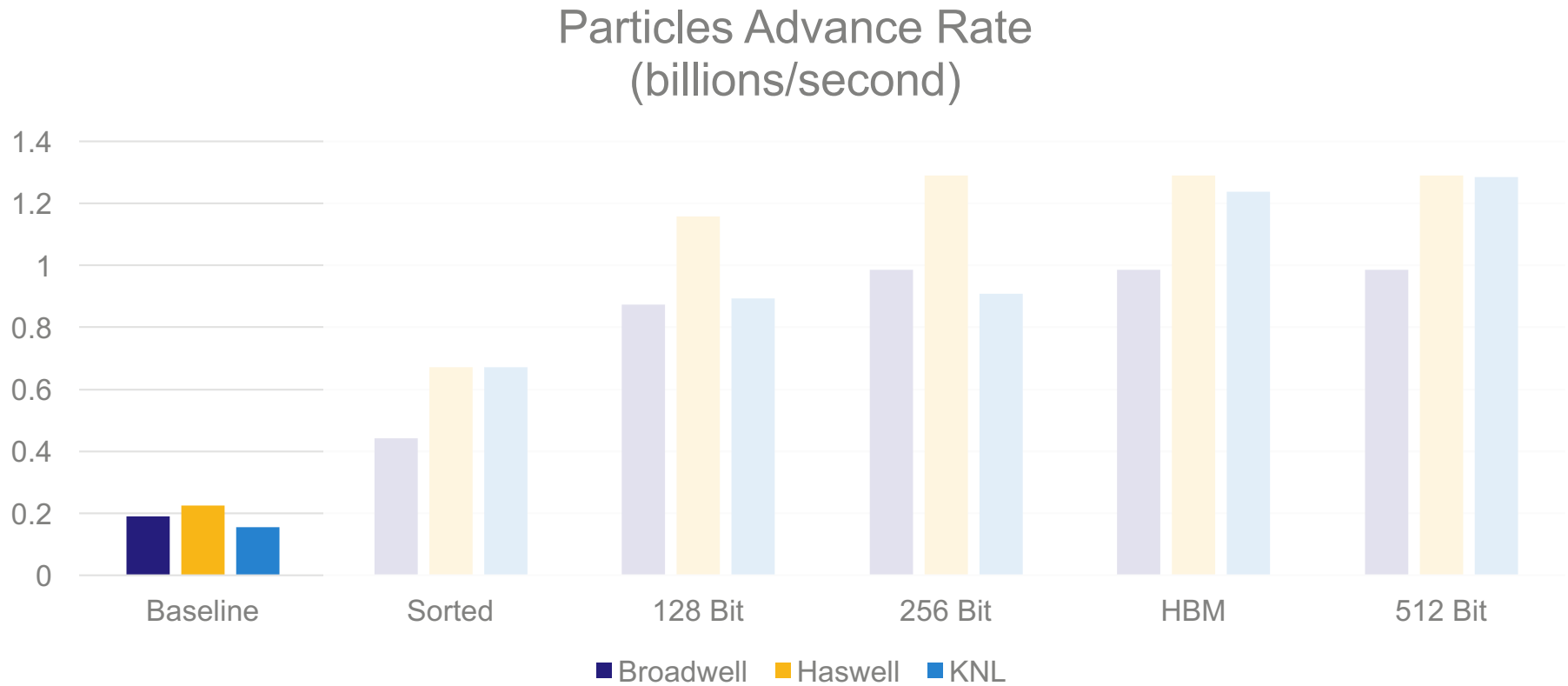- **Explicit short/wide vectorization using hardware intrinsics**

VPIC
Analysis Problems

- **Field Mesh (same for both problems)**
  - 544x96x96 (more cells in the *x*-direction)
  - ~5,000,000 cells
  - Ion and electron species
- **Large Problem (exceeds HBM capacity on KNLs)**
  - 250 particles per cell
  - ~80 GB Memory
- **Small Problem (fits HBM capacity on KNLs)**
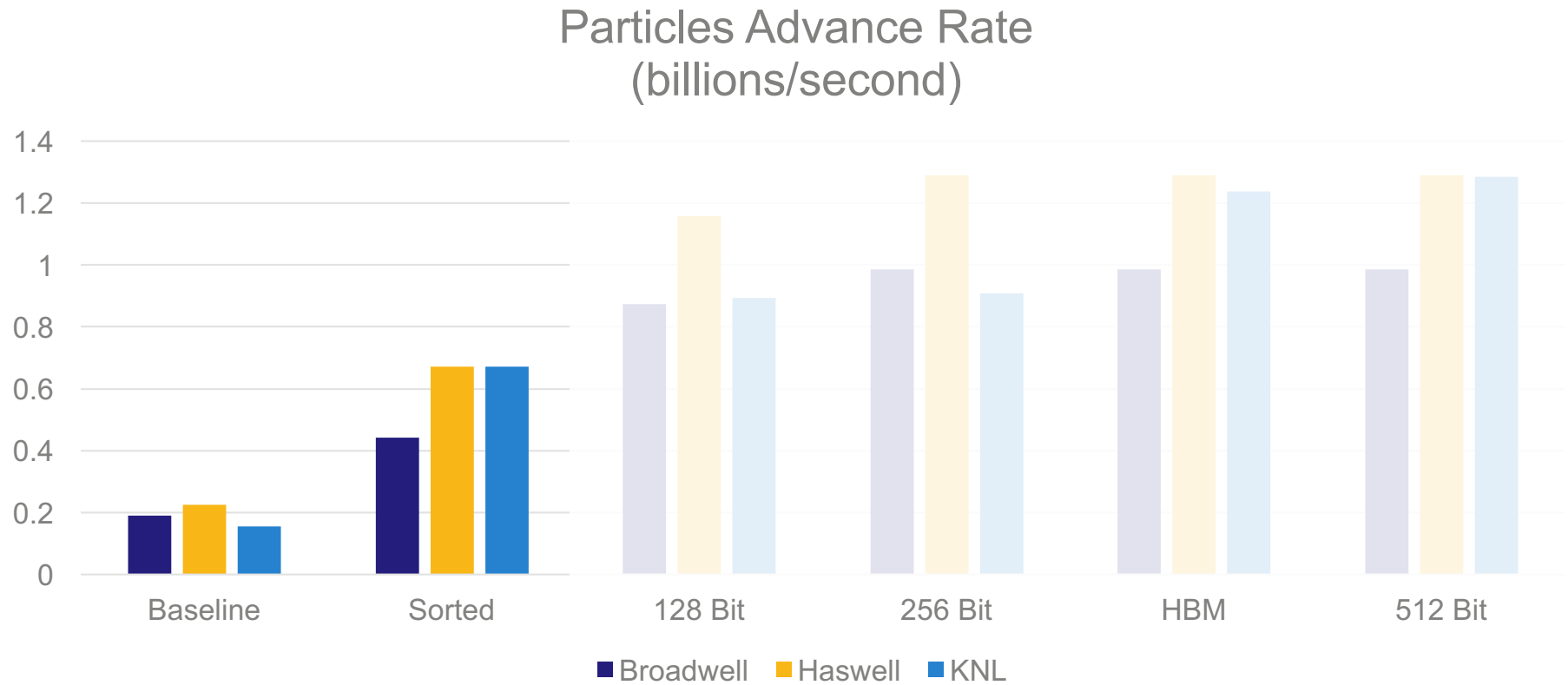 - 25 particles per cell
  - ~8 GB Memory
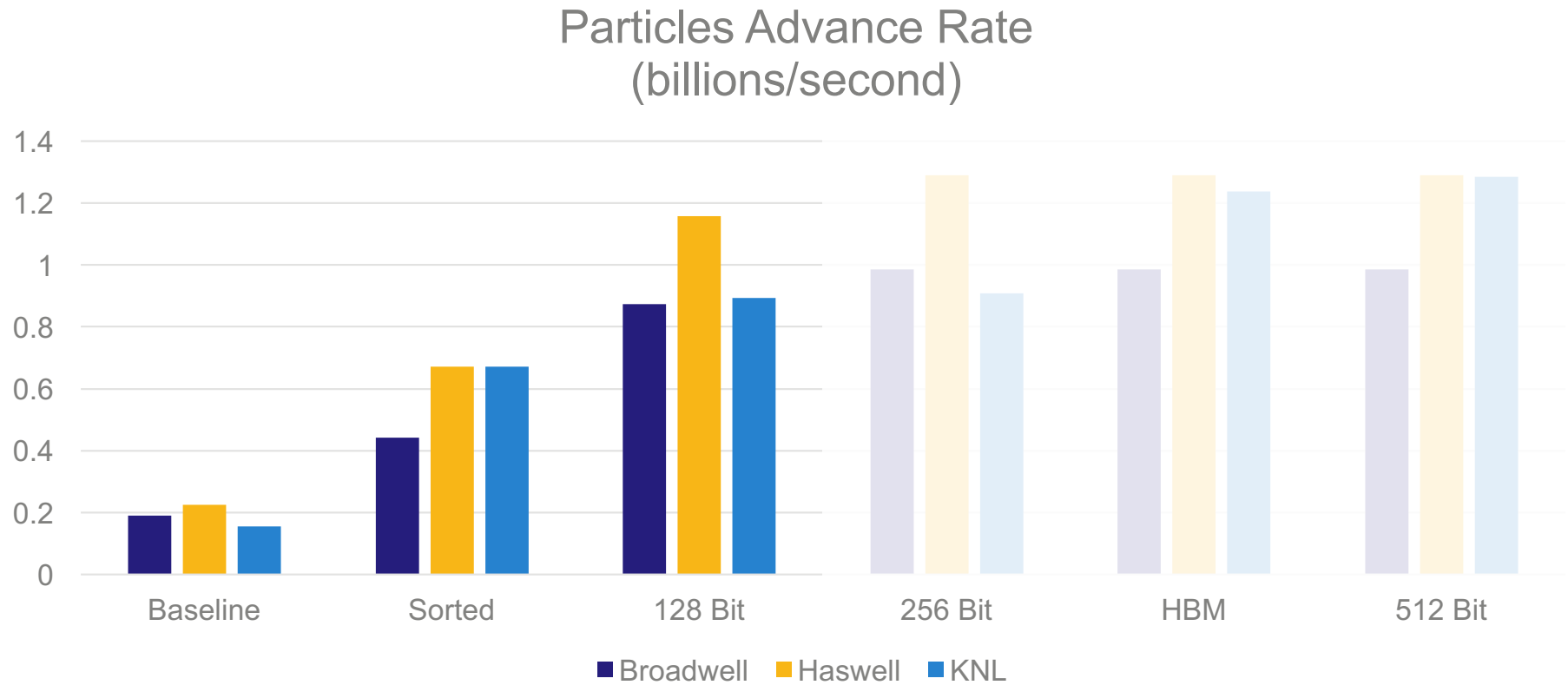
VPIC
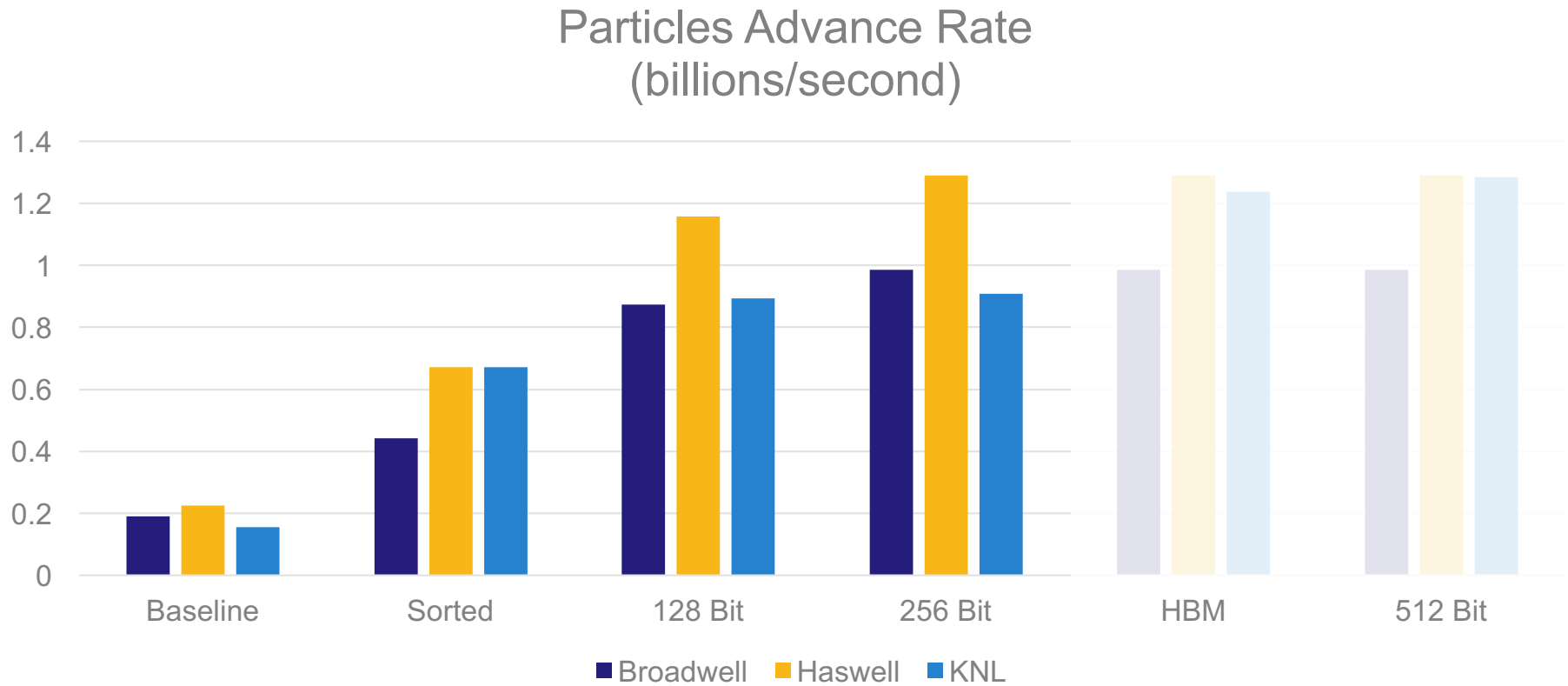# Performance – Large Problem

Particles Advance Rate
(billions/second)

VPIC
Performance – Large Problem

Particles Advance Rate
(billions/second)

VPIC
Performance – Large Problem

Particles Advance Rate
(billions/second)

VPIC
Performance – Large Problem

Particles Advance Rate
(billions/second)

VPIC
Performance – Large Problem

Particles Advance Rate
(billions/second)

VPIC
Performance – Large Problem

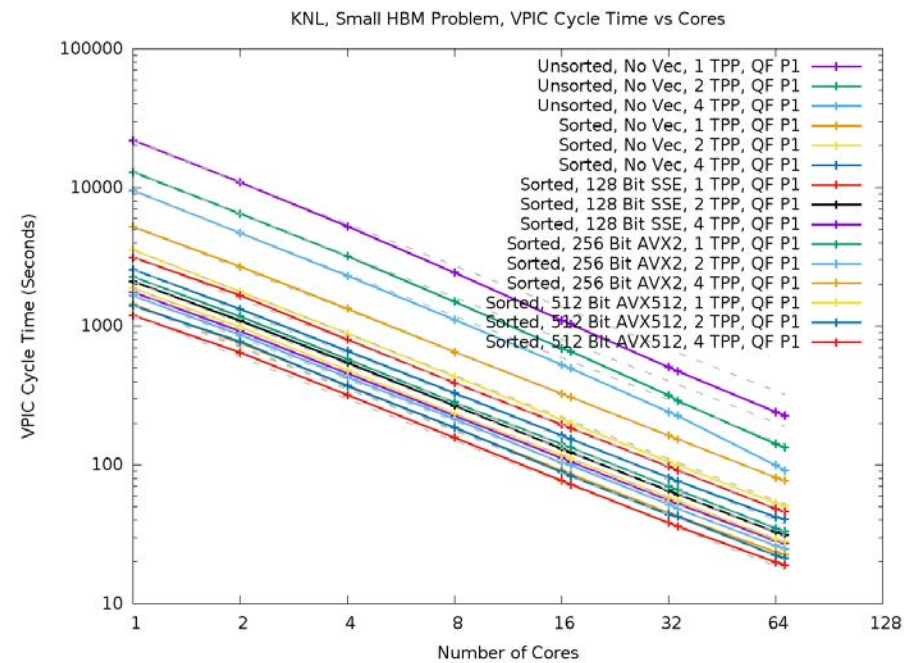

Particles Advance Rate (billions/second)

VPIC
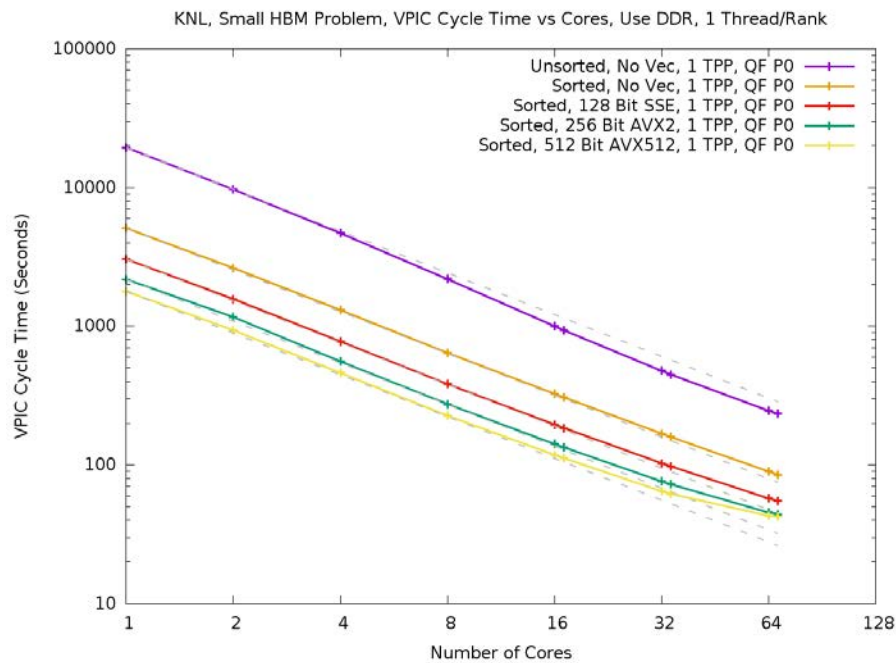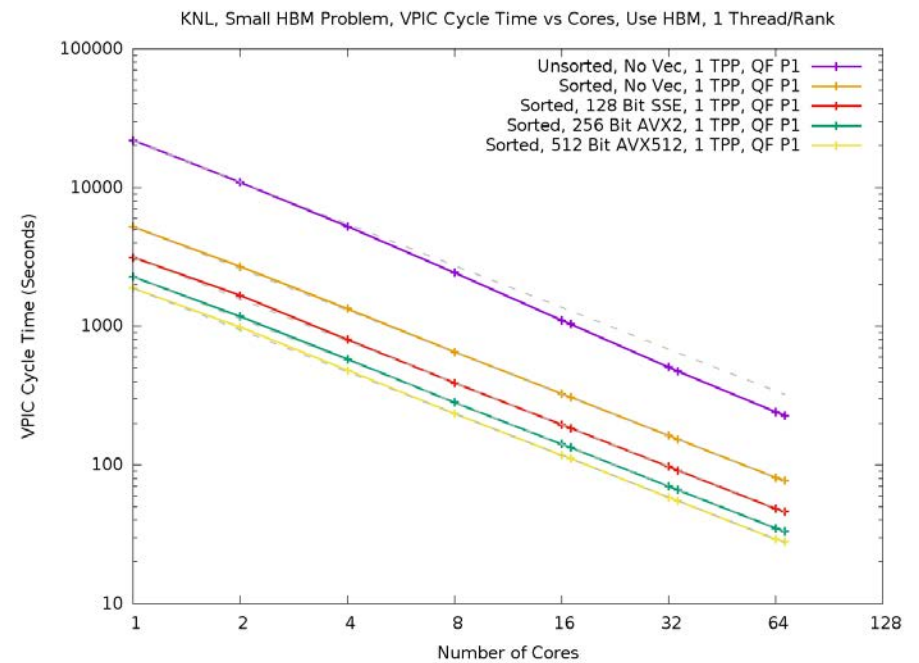# Strong Scalability

Large Problem

Small Problem

VPIC
Memory Bandwidth Limitations – Small Problem

DDR Memory

HBM Memory

VPIC
Memory Bandwidth Limitations – Small Problem

DDR Memory

HBM Memory

VPIC
Memory Bandwidth Limitations – Small Problem

DDR Memory

KNL, Small HBM Problem, VPIC Cycle Time vs Cores, Use DDR, 4 Threads/Rank

Unsorted, No Vec, 4 TPP, QF P0
Sorted, No Vec, 4 TPP, QF P0
Sorted, 128 Bit SSE, 4 TPP, QF P0
Sorted, 256 Bit AVX2, 4 TPP, QF P0
Sorted, 512 Bit AVX512, 4 TPP, QF P0

HBM Memory

KNL, Small HBM Problem, VPIC Cycle Time vs Cores, Use HBM, 4 Threads/Rank

Unsorted, No Vec, 4 TPP, QF P1
Sorted, No Vec, 4 TPP, QF P1
Sorted, 128 Bit SSE, 4 TPP, QF P1
Sorted, 256 Bit AVX2, 4 TPP, QF P1
Sorted, 512 Bit AVX512, 4 TPP, QF P1

VPIC
Memory Bandwidth Limitations – Small Problem

DDR Memory

HBM Memory

HBM is 2x faster

VPIC
Memory Bandwidth Utilization – Large Problem

VPIC
Memory Bandwidth Utilization – Large Problem

DRAM

HBM

more interleaved ➡

sorted

- **Sorted particles essentially get full memory bandwidth utilization**
- **As interleaving occurs, HBM utilization increases due to field data accesses**

VPIC
Memory Bandwidth Utilization – Small Problem

VPIC
Port Utilization

Ports In Use	Cycles Fraction
0	64.7%
1	21.6%
2	15.4%
3+	19.4%

- **For 21.6% of the cycles, the CPU executed only 1 μop on any port**
- **Discussion with Intel suggested that this may be a structural hazard caused by shuffle operations required to re-order data for vector arithmetic**

VPIC
Summary

- **Large problem is currently memory-bandwidth limited**

- **Small problem is limited by instruction latency**

- **We have several strategies to improve HBM bandwidth utilization**

  - Re-implementation of transpose operation to use different intrinsics

  - Try gather/scatter support to re-order data by hand

  - Change data layout: Array-of-Structure-of-Vector (AoSoV)

- **Better HBM utilization may provide an impetus to try triple buffering approach**

FleCSALE
Work done by Marc Charest

FleCSALE
Runtime Portability

FleCSI allows switch between MPI and Legion runtime models with no change to application code

MPI

Legion

Exceptional service in the national interest

Sandia National Laboratories

Metrics for Sandia Applications

S.D. Hammond, C.T. Vaughan, P.T. Lin, D.C. Dinge, C. Hughes, J. Cook, D.M. Pase and R.J. Hoekstra

Application Performance Team (APT), Center for Scientific Computing

Sandia National Laboratories/NM

Outline

- FY17 has been a very busy year for the APT team
  - Analysis of ATDM EMPIRE/Drekar code
  - Analysis of ATDM SPARC code (using native and Trilinos solvers)
  - Support for IC/NALU and ATDM SPARC runs on Trinity Phase-II (KNL) partition
  - Vendor collaborations (see a few examples later in this talk but many more)
  - ASC Advanced Testbed bring up for ARM/Cavium and Intel Sky Lake
  - General ASC/IC (SIERRA & RAMSES) and research code support
  - Profiling and Debug Tools Testing (and feedback to vendors/ISVs)
  - Continued development of LMDS/OVIS Continuous Performance Monitoring

- Slide deck contains some results from many of these activities
- Focus here on a select few "stories" less than specific metrics (see additional slides for more details), <u>huge effort on behalf of the team to cover many areas</u> (thank you!)

Analysis "Stories"

- In today's talk we are going to talk about two codesign studies from this year (based on metrics we gathered):

    - **MPI NIC Resource Requirements** – how big do some significant structures need to be?

    - **Compiler Vectorization Levels** – how do multiple generations of compilers compare and what level of vectorization are we getting?

- We have collected lots more metrics/data on codes and have them in the annotated slides for the review committee

    - See the included slides and let us know if you have any questions

    - Work continues to drill down further on these codes and we will be adding continuous performance monitoring and KokkosP-based overnight testing to our suites in FY18

APPLICATION MPI/NIC METRICS

Codesign of Next-Generation Network Interconnects

MPI Receive Lists (with Cray)

- Working with Cray interconnect team to evaluate *application* requirements for NIC and MPI acceleration
    - MPI Receive Queues (how many pending receives per rank)
    - Average lengths of unexpected queues

- Codesign aspects:
    - Resource structure sizes on the NIC (how many list entries are required)
    - Are our applications well behaved with respect to resource scaling by MPI ranks?

- Implications for future
    - Want to "right size" the resource entries in future network interconnects/NICs

MPI Receive Queue Analysis

NALU MPI Match List Sizes (Per Rank)

- ◆ MPI Only (Max Q Length)
- ▲ MPI Only (Avg, Unexpected)

Efficient use of resources

Drekar MPI Match List Sizes (Per Rank)

- ◆ MPI Only (Max Q Length)
- ■ MPI + OpenMP x 8 (Max Q Length)
- ▲ MPI Only (Avg, Unexpected)
- ■ MPI + OpenMP x 8 (Avg, Unexpected)

Poor scaling with respect to ranks

Move to hybrid helps reduce demands significantly

Conclusions (from Cray)

- The application data has similar behavior to a broad collection of mini-apps
  - If you take behavior of mini-applications across broad cross-section
  - See lots of cases where mini-apps and applications have extra resource usage on rank-0

- Drekar MPI-only will place significant strain on network interconnect/NIC when used at scale
  - Concern this may cause high cost if the number of receives exceeds NIC resources

- See work from Mike Levenhagen, Ryan Grant etc on deep investigation
  - Extension of OpenMPI so we can perform similar studies on CTS and POWER-EA systems
  - Early results also shared with Cray

COMPILER QUALITY/OUTPUT ANALYSIS

Comparing Compiler Technology Generations

Generations of Compilers

- All DOE laboratories receive new/updated compilers at least yearly
    - Improved language support
    - Fixed bugs
    - Better code/instruction generation or new hardware targets
    - Higher level of optimization

- One question is how much better are they getting for our applications in several metrics
    - Time
    - Vectorization
    - Optimization for memory transfers
    - Use of newer hardware functionality

Generations of Compilers

- For Knights Landing these questions come up a lot (especially vectorization)

- Analysis:
  - Study the IC/NALU-CFD Application which is representative of SIERRA STK Mesh use
  - Solves fairly simple "rotor blade" problem using optimized KNL builds
    - Limited OpenMP due to continuing work with Trilinos

- Used Intel 16.2, 17.4 and 18.0 (Beta) compilers
  - *"Compilers are getting better at producing higher levels of vectorization and code optimization over time – true or false?"*

Execution Time

- Intel 18.0 compiler is the fastest (in DDR the results are virtually identical)

- Still see huge differences in threaded versus MPI-only performance because significant number of kernels are not threaded (development on-going)

- Only marginal improvement from HBM/Cache over DDR

NALU R2 Benchmark Problem

Memory Bandwidth Analysis (KNL)

See some improvements in the use of the memory systems but difference is only 2X despite ~5X hardware difference increasing dictated by core performance

Vectorization (Per Rank) (APEX/Tesserae)

Sandia National Laboratories

Approx 15% is FP Instruction

Floating Point Intensity

- Non-FP Ins
- Non-Masked FP Ins
- Masked FP Ins

Approx 20% of FP is vectorized

Vectorization Breakdown

- Scalar FP Ins
- Masked FP Ins
- Non-Masked Vec FP Ins

Approx 35% is "wide" vectorized

Vector Op Breakdown

- 2-Ops
- 4-Ops
- 8-Ops
- 16-Ops

"Wide vectors (AVX2/AVX512)"

Fairly significant improvement from 16.2 to 17.4 in instruction count, breakdown similar

Metrics of Interest (Vendors Ask For)

- **Memory Transferred/Memory Ref**
  - Intel 16.2 – 9.24 Bytes/Ref
  - Intel 17.4 – 9.55 Bytes/Ref
  - Intel 18.0 – 9.72 Bytes/Ref
  - Higher means more efficient transfers

- **Bytes to Floating Point Ratio**
  - Intel 16.2 – 12.86
  - Intel 17.4 – 11.97
  - Intel 18.0 – 11.72
  - Lower means less data movement

Source/Destination for Calculations

Legend: ■ Reg/Reg ■ Reg/Mem ■ VecReg/VecReg ■ VecReg/Mem

Y-axis: Instructions (0 to 1.4E+13)

X-axis: Intel 16.2, Intel 17.4, Intel 18.0

Reduce data motion/improve execution speed

Fairly significant improvement from 16.2 to 17.4 in instruction counts, gradually more data using vector registers

Discussion

- Seeing improvements in time with successive compiler generations over the same code base
    - Shows there is value in pursuing updating compiler versions (even though revalidation is time consuming)
    - Allows us to more efficiently use hardware resources which we have paid for

- Vectorization levels are still lower than we would like even on a system like Knights Landing where the VPUs are very efficient
    - C++ continues to be our nemesis with vectorization
    - Kokkos adding to the complexity with inlining and use of views that cannot pass alias-analysis when used in lambda functions

- Future hardware systems with greater memory bandwidth <u>are likely to place much greater emphasis on instruction efficiency/optimization of computation</u>

FY17 CODESIGN UPDATE FROM SANDIA

Codesign Highlights for FY17

- Continue to bring up of ATDM codes on Haswell, Knights, POWER8 and GPUs
  - Seeing significant efforts with underpinning of Trilinos and some really big improvements now underway
- Exceptionally broad testing of compilers including LLVM, GCC, IBM-XL, NVCC, Intel, PGI and in some limited cases Cray
  - Early Intel 18.0 working very well for code portfolio
  - 30% reduction in compile times with IBM XL (still takes a long time)
  - Submitted big pile of bugs this year on early releases
- AMD has supplied a ROCm back-end developed for Kokkos (and lots of fixes)
  - Benchmarking on Fiji GPUs underway

https://github.com/kokkos/kokkos

Codesign Highlights for FY17

- Local team performed early runs on Sky Lake Platinum Xeon
  - Results are looking very strong for many benchmark areas
- Wave-1 shake-down of Cavium ARM processors completed
  - Have NALU-CFD and large suite of packages running (even with the correct answer!), results provided back to HPE
- Continuous Performance Monitoring LDMS System (Jim Brandt, Jeanine Cook and Tom Tucker)
  - Slowly taking shape, prototypes on Mutrino HSW/KNL now underway and data gathering shown when used in DAT modes
  - Hoping for some of this to come online in FY18

Speedup Relative to Dual-Socket Haswell

Legend: ■ KNL-DDR ■ KNL-HBM □ Haswell ■ Sky Lake ■ ThunderX2

Y-axis: Speedup (X), from 0 to 5

Categories: STREAM, LULESH, PENNANT, MiniFE

Thank You!

- Sandia Application Performance Team ("APT")
- ASC Advanced Test Bed Admin/Support Team
- SIERRA, RAMSES and CTH Code Groups
- Trilinos Developers
- ATDM Code Groups
- ATDM Software Environment Team

- L2 Milestone Review Team!

The FY17 Award for Patience Goes To..

- Clay Hughes

- For services to getting the IBM XL compiler to compile Trilinos and working with the NVIDIA Profiler despite it taking a significant amount of time

APPLICATION METRIC GATHERING

Application Metric Gathering

- ATDM/SPARC
  - MPI Operation Breakdown
  - MPI Function Times
  - Thread/Rank Scaling
  - Native counter timing comparison (Haswell vs. Knights Landing)
  - Comparison of native solvers vs. Trilinos

- ATDM/EMPIRE/Drekar
  - Native counter timing comparison of Haswell, Knights Landing and POWER w/ P100 GPUs
  - GPU/CPU kernel profiling

- IC/NALU
  - Native counter
  - MPI Match Lists (early in presentation)

Performance Analysis of ATDM/EMPIRE/Drekar

Work by Paul Lin

EMPIRE/Drekar Solver Performance

- Focus was on trying to improve performance of ATDM/EMPIRE/Drekar/fluids solve (Trilinos solvers)
- What application developers (e.g. Drekar dev) really want
  - profile code performance, e.g. provide breakdown where time is being spent with the goal of identifying performance bottlenecks
  - even better: help/work with them or TPL developers to improve performance
- Performance profiling and improving thread scaling on KNL
  - mainly focused on single KNL (mutrino/tr2, ellis/bowman) because if doesn't thread scale on single KNL, will not on multiple KNLs
  - after get good single KNL performance, then will focus on multiple KNLs
  - But did look at performance of multiple KNLs (up to 512)
  - EMPIRE lead provided representative plasma test case
- Performance profiling for large-scale simulations
  - Focus on MPI performance
  - Most of this work was for CPUs rather than KNL
  - Employed a different test case (MHD)

EMPIRE/Drekar Thread Scaling on KNL

- Accomplishments
  - Tpetra matrix-matrix addition (worst thread scaling kernel)
    - Significant reduction in time and significant improvement in thread scaling for for single KNL with one MPI process
    - Still in progress: multiple MPI processes and multiple KNLs
  - Gauss-Seidel (second worst thread scaling kernel)
    - Significant improvement in thread scaling for solve (apply GS)
    - KokkosKernels team working to improve setup time for GS
- Had identified additional major kernels that did not thread scale (work still in progress)
  - few items in MueLu setup
  - Tracking all above issues with github issue

EMPIRE/Drekar Solvers on KNL

- Single mutrino KNL, 1 MPI, increase OMP threads from 1 to 64

- Setup: most expensive kernel for 1 OMP is mat-mat multiply, but KKMEM thread scales well (1 OMP ~47% of setup, 64 OMP ~4% of setup)

- Worst thread scaling kernels:
  1. Tpetra matrix-matrix addition during setup (separate from multigrid)
  2. Application of Gauss-Seidel smoother during solve
  3. A few kernels in MueLu multigrid setup
     - For 1 and 64 OMP, multigrid setup ~9% and ~13% of total setup, respectively

Drekar plasma_perf_test (40^3 elem) single KNL thread scaling original MatMat::add, old Gauss-Seidel

On-Node Improvement (EMPIRE/Drekar)

- Single mutrino KNL, 1 MPI, increase OMP threads from 1 to 64

- Significant reduction in time and significant improvement in thread scaling for for single KNL with one MPI process

- Still in progress: multiple MPI processes and multiple KNLs

Drekar plasma_perf_test (40^3 elem) single KNL thread scaling

Improv. to Gauss Seidel Setup/Apply

- Single mutrino KNL, 1 MPI, increase OMP threads from 1 to 64

- previously used non-threaded Gauss-Seidel ("old GS")

- switched to KokkosKernels threaded Gauss-Seidel ("orig MT-GS")
  - initially, threaded Gauss-Seidel setup time was very high, worked with KokkosKernels team (Deveci) to significantly reduce

- Improved MT Gauss-Seidel ("new MT-GS"): good threaded performance for setup to 8 OMP, reasonable threaded solve
  - KokkosKernels team currently working on more improvements

Drekar plasma_perf_test (40^3 elem) single KNL thread scaling without improved MatMat::add; compare diffierent Gauss-Seidels

- old GS: setup+solve
- orig MT-GS: setup+solve
- new MT-GS: setup+solve

EMPIRE/Drekar Large Scale Runs

- Focused on MPI performance and scaling rather than threaded
- Accomplishments:
  - panzer scaling bug after change from stk_classic to stk (identified and now fixed)
    - Caused a major memory explosion, Drekar was failing with 32k MPI processes when it had previously been run with 1.6M MPI
  - STK CommSparse scaling bug (identified and fixed)
    - Big performance issue at scale; both memory and time explosion
      - same STK bug and Cray MPI bug that caused Nalu to fail the Trinity Phase 2 acceptance test; I was actually simultaneously chasing down this issue with both Drekar and Nalu, not realizing at that time it was the same bugs
  - MueLu setup
    - identified and work with MueLu team to fix a modest scaling issue, for 524,288 MPI reduced setup by ~10%
    - we worked with Zoltan2 team (M. Deveci) on a scaling issue---preliminary results on 131,072 MPI on BG/Q showed a reduction in setup of ~30%
- Identified other issues (work still in progress)
  - MueLu setup
    - Few other MueLu setup scaling issues (working with MueLu team to resolve)
    - Tpetra matrix-matrix multiplication
  - Scaling issues with periodic BCs (panzer scaling issue; pushed back to FY18)

EMPIRE/DREKAR ANALYSIS ON POWER-EA

Work by Clay Hughes
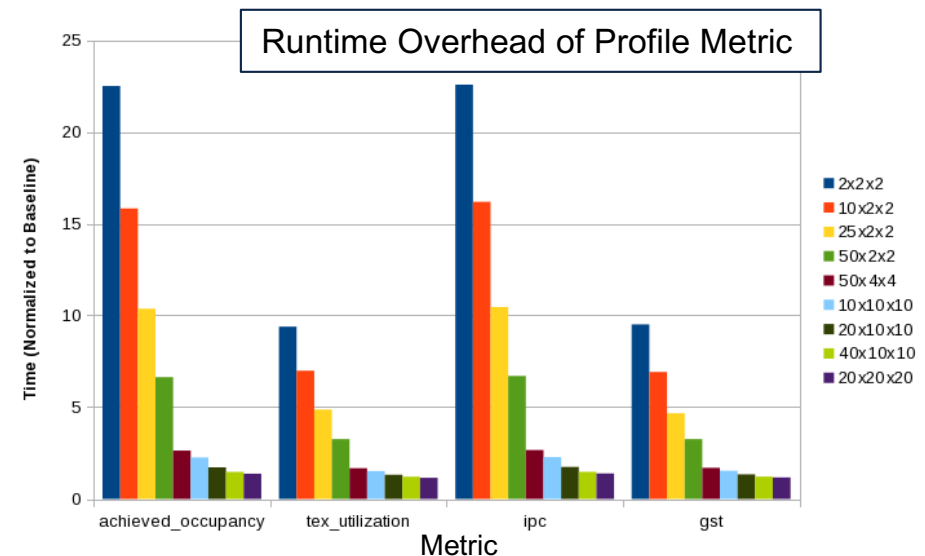
EMPIRE/Drekar on POWER-EA

- GCC + CUDA Builds working successfully although problem sizes often need to be reduced to fit into GPU memory
  - Finding that scratch arrays used in inner kernels are growing as kernel complexity increases

- Significant amount of work trying to get IBM XL to compile code so we can perform a cross-comparison (performance and functionality)
  - Latest IBM XL 14.1.0 does not work correctly
  - Either does not compile the code (errors in particular with OpenBLAS inline assembly for VSX)
  - Or compiles and has a segmentation fault in BLAS routines (not present in GCC version of the code)

- Metric analysis of EMPIRE/Drekar on Pascal GPUs an issue

Performing Metric Gathering on P8-EA

- Gathering metrics using NVIDIA profiler in POWER-EA systems is causing significant increases in runtime

- Specific metrics causing problems:
  - DRAM Read/Write Throughput
  - L2 Read/Write Throughout
  - GLD

- In some cases the runtimes exceed 48 hours which are our job queue limits
  - The worst profiling overhead we have seen in tools

Runtime Overhead of Profile Metric

Outcomes

- **Profiling overhead bug is reported to NVIDIA**
  - Bug: #1977033, current status is open, NVIDIA need small reproducer
  - Not sure that small reproducer will show the outcome so we are looking at whether SDK samples will induce the behavior

- **Potential IBM XL Inline Assembly Bug in XL 14.1.0 Beta**
  - Reported to IBM Toronto compiler team
  - Investigation underway

- **Feedback to EMPIRE/Drekar development on the use of large "work sets" which can significantly bloat memory consumption on GPU**
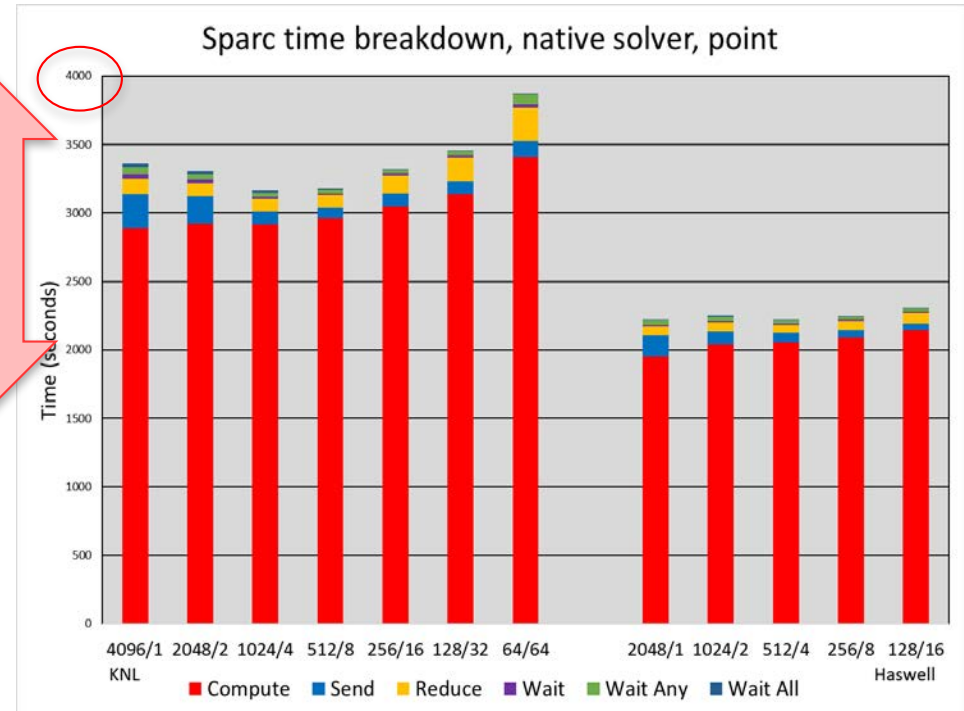
ANALYSIS OF ATDM/SPARC

Work by Courtenay Vaughan

SPARC Analysis

- Performed for the Trinity Phase-II open-campaign to assess scaling

- SPARC is written to utilize two approaches:
  - Native Solver written in OpenMP as prototype (considered to be very fast)
  - Full Trilinos solvers being slowly ported to Kokkos (many packages, slower but being optim'd)

- Questions:
  - What needs most improvement?
  - What packages?
  - How is scaling on and between nodes performing today?

SPARC Runtime with Solver Variants

Up to 2X performance difference (native faster) but .. MPI times are slightly faster in Trilinos (shows history as *MPI-only code*, SPARC shows what a *newly* threaded, from group up code *can do*.

Thread Scaling of SPARC on KNL

- Captured with in-code timings but also checked against CrayPAT outputs (for small nodes)

- Trilinos Line Solver is approx. 50% slower

- Trilinos Point Solver is competitive but does not scale as well with large thread counts

Observations from MPI Analysis

- Number of messages exchanged decreases from 130000 to 103000 per MPI rank when going from MPI only to 64 threads
  - So the total number of messages decreases substantially
- The bytes exchanged grows from 1.2 E10 to 1.46 E11 bytes per MPI rank when going from MPI only to 64 threads
  - The total message traffic drops from 4.5 E13 to 9.33 E12 bytes
  - The average message size goes from 82 KB to 1.3 MB
- For the native solver, the number of reduces stays constant at 47500 calls and 3.6 MB
- For the Trilinos solver, the number of reduces stays constant at 30000 calls and 2.4 MB

Update from SPARC Team

- Andrew Bradley has been developing a faster solver for Trilinos (using experiments from native solver)
  - Early results which were used for the Power API L2 show significant improvement
  - Now as fast/faster than the native solver
  - Builds on experiences from native development but with additional optimizations

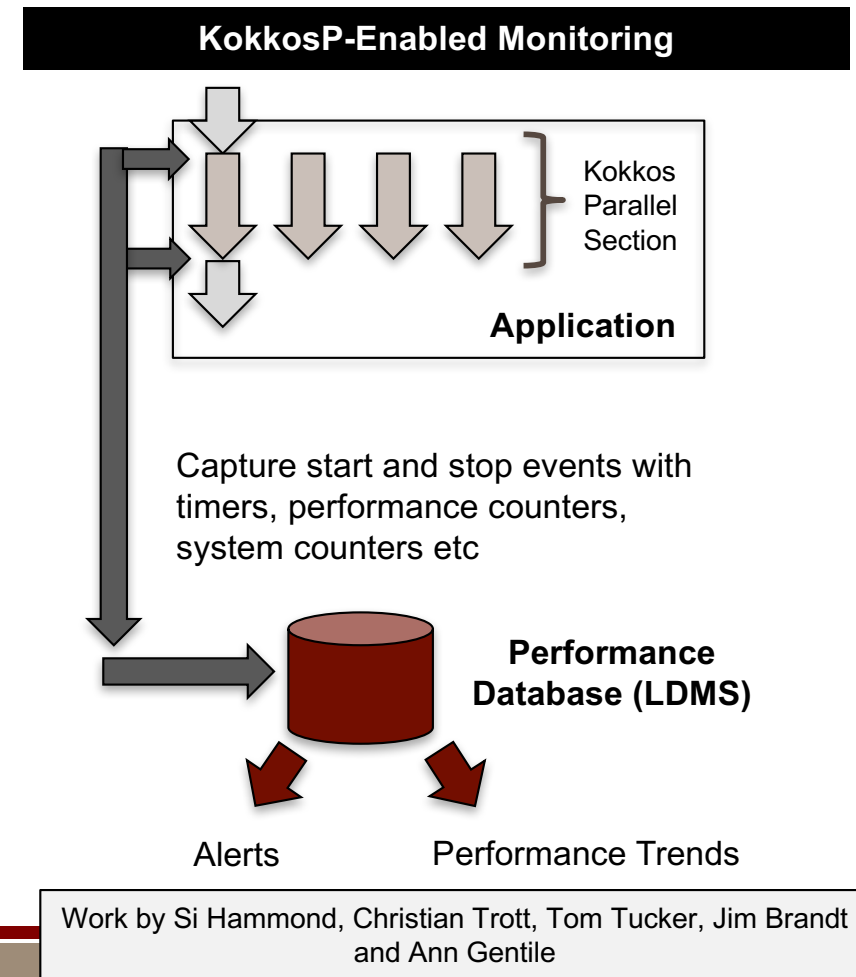- Will be pushed into Trilinos before the end of the FY

LMDS Kernel and Continuous Performance Tracing

Kernel Monitoring Work by Si Hammond, Christian Trott, Tom Tucker, Jim Brandt and Ann Gentile

Continuous Monitoring Work by Jeanine Cook, Jim Brandt, Ann Gentile and Tom Tucker

Application Kernel Monitoring

- **KokkosP-based kernel profiling direct into LDMS metrics database**
  - Demonstrated on KNL, Haswell and POWER8
  - Ability to hook into Kokkos kernels and regions without requiring changes to compiler, linking etc

- **Provides detailed drill down for Kokkos-enabled applications**
  - Can connect with vendor tools, PAPI, perf, etc
  - Data movement (e.g. NVLINK, MC-DRAM etc)
  - Provides Kokkos context (e.g. "parallel-for")

KokkosP-Enabled Monitoring

Kokkos Parallel Section

Application

Capture start and stop events with timers, performance counters, system counters etc

Performance Database (LDMS)

Alerts   Performance Trends

Work by Si Hammond, Christian Trott, Tom Tucker, Jim Brandt and Ann Gentile

LDMS Continuous Monitoring

- Continuous HPC system monitoring framework that:
    - Collects numeric data
    - Moves and aggregates data
    - Stores data
    - Analyze data
        - Troubleshooting
        - Optimization
        - Inform future systems
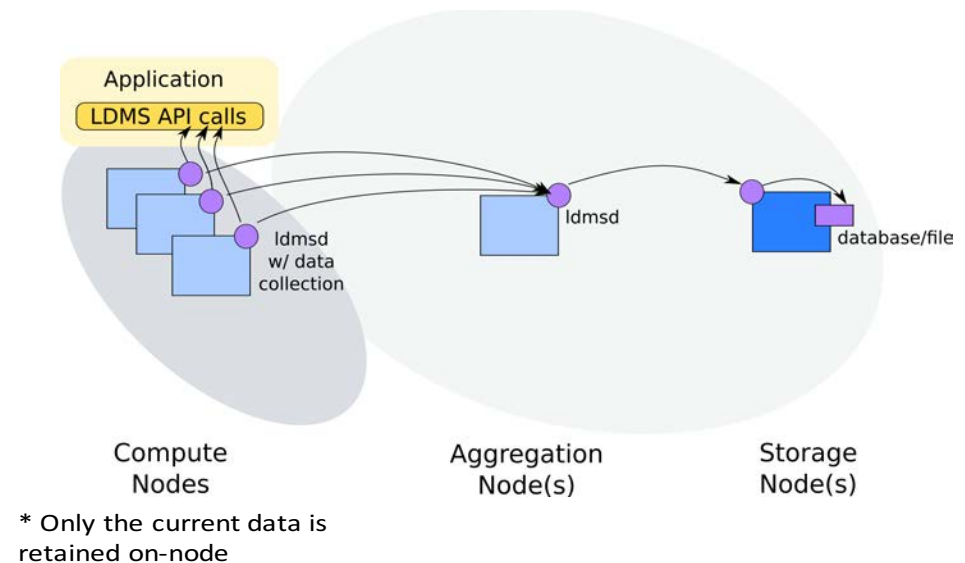- Plug-in architecture
    - Customized samplers for data collection
- Backend
    - Storage (SOS DB)
    - Analysis (Python)
    - Visualization (Grafana)
- Application-based samplers
    - PAPI, Perf, MPI, DLoad

Application
LDMS API calls

ldmsd w/ data collection

ldmsd

database/file

Compute Nodes

Aggregation Node(s)
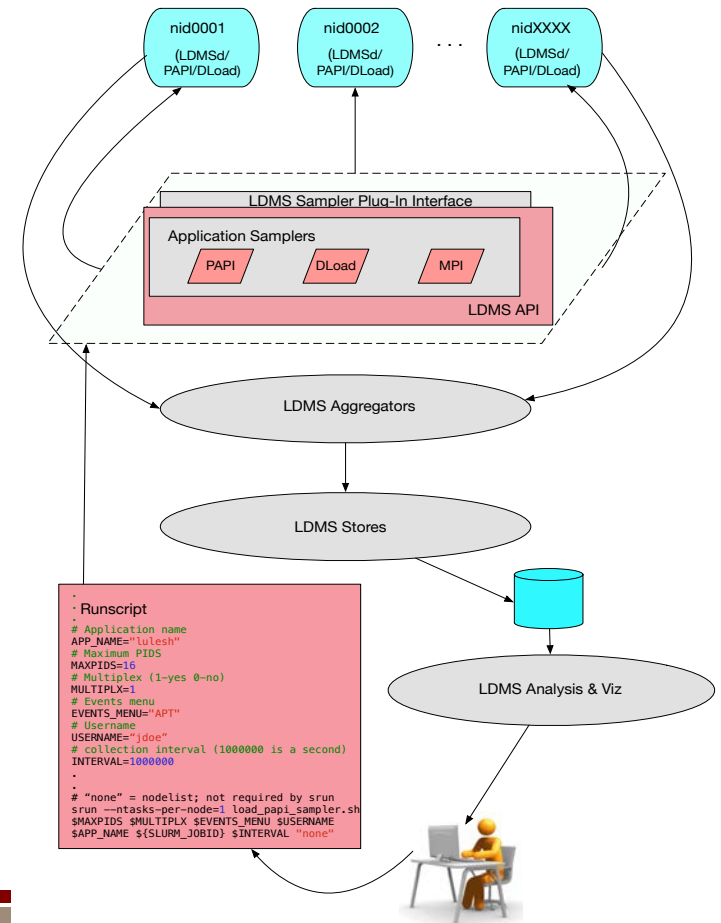
Storage Node(s)

* Only the current data is retained on-node

Work by Jeanine Cook, Tom Tucker, Jim Brandt and Ann Gentile

Application LDMS "Samplers"

- PAPI
    - Menu-driven
        - Instruction mix (vectorization, cycles per instruction (CPI))
        - Branch mispredict rates
        - Cache/memory statistics (miss/hit rates)
        - Average (estimated) memory bandwidth
        - Stall cycle profiles (which components incur stalls)
        - Other events for performance characterization work
    - Cross-platform
        - Haswell/Broadwell, Ivybridge
        - Very near future: Power8, KNL
    - Opt-in and configure through runscript variables
- DLoad
    - Enables dynamic loading and reconfiguring (e.g., application/PAPI events) of application samplers
- MPI
    - Developed by UCF
    - Working on usability issues
- Perf
    - Complete after PAPI development is complete and tested
    - Currently low priority

LDMS Status

- PAPI & DLoad samplers
  - Small-scale testing complete on several testbeds
  - Large-scale testing on Mutrino/Haswell scheduled 08/31
  - All KNL testing by 11/01

- MPI sampler
  - Targeting testing completion by Sept 15

- LDMS w/Application Samplers, and Backend, Analysis, basic Viz
  - Plan to launch on Mutrino/Haswell by Sept. 15

- Working with Sierra team to integrate LDMS application monitoring into nightly testing
  - Hope to complete integration and testing by 12/01